
Improving the Expected Improvement Algorithm

Anonymous Author(s)

Affiliation

Address

email

Abstract

1 The expected improvement (EI) algorithm is a popular strategy for information
2 collection in optimization under uncertainty. The algorithm is widely known to
3 be too greedy, but nevertheless enjoys wide use due to its simplicity and ability
4 to handle uncertainty and noise in a coherent decision theoretic framework. To
5 provide rigorous insight into EI, we study its properties in a simple setting of
6 Bayesian optimization where the domain consists of a finite grid of points. This
7 is the so-called best-arm identification problem, where the goal is to allocate
8 measurement effort wisely to confidently identify the best arm using a small
9 number of measurements. In this framework, one can show formally that EI is far
10 from optimal. To overcome this shortcoming, we introduce a simple modification
11 of the expected improvement algorithm. Surprisingly, this simple change results in
12 an algorithm that is asymptotically optimal for Gaussian best-arm identification
13 problems, and provably outperforms standard EI by an order of magnitude.

14 1 Introduction

15 Recently Bayesian optimization has received much attention in the machine learning community
16 [23]. This literature studies the problem of maximizing an unknown black-box objective function by
17 collecting noisy measurements of the function at carefully chosen sample points. At first a prior belief
18 over the objective function is prescribed, and then the statistical model is refined sequentially as data
19 are observed. *Expected improvement (EI)* [14] is one of the most widely-used Bayesian optimization
20 algorithms. It is a greedy improvement-based heuristic that samples the point offering greatest
21 expected improvement over the current best sampled point. EI is simple and readily implementable,
22 and it offers reasonable performance in practice.

23 Although EI is reasonably effective, it is too greedy, focusing nearly all sampling effort near the
24 estimated optimum and gathering too little information about other regions in the domain. This
25 phenomenon is most transparent in the simplest setting of Bayesian optimization where the function's
26 domain is a finite grid of points. This is the problem of best-arm identification (BAI) [2] in a multi-
27 armed bandit. The player sequentially selects arms to measure and observes noisy reward samples
28 with the hope that a small number of measurements enable a confident identification of the best
29 arm. Recently Ryzhov [22] studied the performance of EI in this setting. His work focuses on a link
30 between EI and another algorithm known as the optimal computing budget allocation [4], but his
31 analysis reveals EI allocates a vanishing proportion of samples to suboptimal arms as the total number
32 of samples grows. Any method with this property will be far from optimal in BAI problems [2].

33 In this paper, we improve the EI algorithm dramatically through a simple modification. The resulting
34 algorithm, which we call *top-two expected improvement (TTEI)*, combines the top-two sampling
35 idea of Russo [21] with a careful change to the improvement-measure used by EI. We show that
36 this simple variant of EI achieves strong asymptotic optimality properties in the BAI problem, and
37 benchmark the algorithm in simulation experiments.

38 Our main theoretical contribution is a complete characterization of the asymptotic proportion of
39 samples TTEI allocates to each arm as a function of the true (unknown) arm means. These particular
40 sampling proportions have been shown to be optimal from several perspectives [5, 13, 10, 21, 9], and
41 this enables us to establish two different optimality results for TTEI. The first concerns the rate at
42 which the algorithm gains confidence about the identity of the optimal arm as the total number of
43 samples collected grows. Next we study the so-called fixed confidence setting, where the algorithm is
44 able to stop at any point and return an estimate of the optimal arm. We show that when applied with
45 the stopping rule of Garivier and Kaufmann [9], TTEI essentially minimizes the expected number of
46 samples required among all rules obeying a constraint on the probability of incorrect selection.

47 One undesirable feature of our algorithm is its dependence on a tuning parameter. Our theoretical
48 results precisely show the impact of this parameter, and reveal a surprising degree of robustness to its
49 value. It is also easy to design methods that adapt this parameter over time to the optimal value, and
50 we explore one such method in simulation. Still, removing this tuning parameter is an interesting
51 direction for future research.

52 **Further related literature.** Despite the popularity of EI, its theoretical properties are not well
53 studied. A notable exception is the work of Bull [3], who studies a global optimization problem and
54 provides a convergence rate for EI's expected loss. However, it is assumed that the observations
55 are noiseless. Our work also relates to a large number of recent machine learning papers that try to
56 characterize the sample complexity of the best-arm identification problem [6, 19, 2, 8, 15, 11, 12, 16–
57 18]. Despite substantial progress, matching asymptotic upper and lower bounds remained elusive in
58 this line of work. Building on older work in statistics [5, 13] and simulation optimization [10], recent
59 work of Garivier and Kaufmann [9] and Russo [21] characterized the optimal sampling proportions.
60 Two notions of asymptotic optimality are established: *sample complexity in the fixed confidence*
61 *setting* and *rate of posterior convergence*. Garivier and Kaufmann [9] developed two sampling
62 rules designed to closely track the asymptotic optimal proportions and showed that, when combined
63 with a stopping rule motivated by Chernoff [5], this sampling rule minimizes the expected number
64 of samples required to guarantee a vanishing threshold on the probability of incorrect selection is
65 satisfied. Russo [21] independently proposed three simple Bayesian algorithms, and proved that
66 each algorithm attains the optimal rate of posterior convergence. TTEI proposed in this paper is
67 conceptually most similar to the top-two value sampling of Russo [21], but it is more computationally
68 efficient.

69 1.1 Main Contributions

70 As discussed below, our work makes both theoretical and algorithmic contributions.

71 **Theoretical:** Our main theoretical contribution is Theorem 1, which establishes that TTEI—a simple
72 modification to a popular Bayesian heuristic—converges to the known optimal asymptotic
73 sampling proportions. It is worth emphasizing that, unlike recent results for other top-two
74 sampling algorithms [21], this theorem establishes that the expected time to converge to the
75 optimal proportions is finite, which we need to establish optimality in the fixed confidence
76 setting. Proving this result required substantial technical innovations. Theorems 2 and 3
77 are additional theoretical contributions. These mirror results in [21] and [9], but we extract
78 minimal conditions on sampling rules that are sufficient to guarantee the two notions of
79 optimality studied in these papers.

80 **Algorithmic:** On the algorithmic side, we substantially improve a widely used algorithm. TTEI can
81 be easily implemented by modifying existing EI code, but, as shown in our experiments, can
82 offer an order of magnitude improvement. A more subtle point involves the advantages of
83 TTEI over algorithms that are designed to directly target convergence on the asymptotically
84 optimal proportions. In the experiments, we show that TTEI substantially *outperforms an*
85 *oracle sampling rule* whose sampling proportions directly track the asymptotically optimal
86 proportions. This phenomenon should be explored further in future work, but suggests that
87 by carefully reasoning about the value of information TTEI accounts for important factors
88 that are washed out in asymptotic analysis. Finally—as discussed in the conclusion—although
89 we focus on uncorrelated priors we believe our method can be easily extended to more
90 complicated problems like that of best-arm identification in linear bandits [24].

91 **2 Problem Formulation**

92 Let $A = \{1, \dots, k\}$ be the set of arms. The reward $Y_{n,i}$ of arm $i \in A$ at time $n \in \mathbb{N}$ follows a
 93 normal distribution $N(\mu_i, \sigma^2)$ with common known variance σ^2 , but unknown mean μ_i . At each
 94 time $n = 1, 2, \dots$, an arm $I_n \in A$ is measured, and the corresponding noisy reward Y_{n,I_n} is observed.
 95 The objective is to allocate measurement effort wisely in order to confidently identify the arm with
 96 highest mean using a small number of measurements. We assume that $\mu_1 > \mu_2 > \dots > \mu_k$, i.e., the
 97 arm-means are unique and arm 1 is the best arm. Our analysis takes place in a *frequentist setting*, in
 98 which the true means (μ_1, \dots, μ_k) are fixed but unknown. The algorithms we study, however, are
 99 Bayesian, in the sense that they begin with prior over the arm means and update the belief to form a
 100 posterior distribution as evidence is gathered.

101 **Prior and Posterior Distributions.** The sampling rules studied in this paper begin with a normally
 102 distributed prior over the true mean of each arm $i \in A$ denoted by $N(\mu_{1,i}, \sigma_{1,i}^2)$, and update this to
 103 form a posterior distribution as observations are gathered. By conjugacy, the posterior distribution
 104 after observing the sequence $(I_1, Y_{1,I_1}, \dots, I_{n-1}, Y_{n-1,I_{n-1}})$ is also a normal distribution denoted
 105 by $N(\mu_{n,i}, \sigma_{n,i}^2)$. The posterior mean and variance can be calculated using the following recursive
 106 equations:

$$\mu_{n+1,i} = \begin{cases} (\sigma_{n,i}^{-2} \mu_{n,i} + \sigma^{-2} Y_{n,i}) / (\sigma_{n,i}^{-2} + \sigma^{-2}) & \text{if } I_n = i, \\ \mu_{n,i}, & \text{if } I_n \neq i, \end{cases}$$

107 and

$$\sigma_{n+1,i}^2 = \begin{cases} 1 / (\sigma_{n,i}^{-2} + \sigma^{-2}) & \text{if } I_n = i, \\ \sigma_{n,i}^2, & \text{if } I_n \neq i. \end{cases}$$

108 We denote the posterior distribution over the vector of arm means by

$$\Pi_n = N(\mu_{n,1}, \sigma_{n,1}^2) \otimes N(\mu_{n,2}, \sigma_{n,2}^2) \otimes \dots \otimes N(\mu_{n,k}, \sigma_{n,k}^2)$$

109 and let $\theta = (\theta_1, \dots, \theta_k)$. For example, with this notation

$$\mathbb{E}_{\theta \sim \Pi_n} \left[\sum_{i \in A} \theta_i \right] = \sum_{i \in A} \mu_{n,i}.$$

110 The posterior probability assigned to the event that arm i is optimal is

$$\alpha_{n,i} \triangleq \mathbb{P}_{\theta \sim \Pi_n} \left(\theta_i > \max_{j \neq i} \theta_j \right). \quad (1)$$

111 To avoid confusion, we use $\theta = (\theta_1, \dots, \theta_k)$ to denote a random vector of arm means drawn from
 112 the algorithm's posterior Π_n , and $\mu = (\mu_1, \dots, \mu_k)$ to denote the vector of true arm means.

113 **Two notions of asymptotic optimality.** Our first notion of optimality relates to the rate of poste-
 114 rior convergence. As the number of observations grows, one hopes that the posterior distribution
 115 definitively identifies the true best arm, in the sense that the posterior probability $1 - \alpha_{n,1}$ assigned
 116 by the event that a different arm is optimal tends to zero. By sampling the arms intelligently, we hope
 117 this probability can be driven to zero as rapidly as possible. We will see that under TTEI the posterior
 118 probability tends to zero at an exponential rate, and so following Russo [21], we aim to maximize the
 119 exponent governing the rate of decay, effectively solving the optimization problem

$$\min_{\text{sampling rules}} \limsup_{n \rightarrow \infty} \frac{1}{n} \log(1 - \alpha_{n,1}).$$

120 The second setting we consider is often called the “fixed confidence” setting. Here, the agent is
 121 allowed at any point to stop gathering samples and return an estimate of the identity of the optimal. In
 122 addition to the sampling rule TTEI, we require a stopping rule that selects a time τ at which to stop,
 123 and decision rule that returns an estimate \hat{i}_τ of the optimal arm based on the first τ observations. We
 124 consider minimizing the average number of observations $\mathbb{E}[\tau]$ required by an algorithm guaranteeing
 125 a vanishing probability δ of incorrect identification, i.e., $\mathbb{P}(\hat{i}_\tau \neq 1) \leq \delta$. Following Garivier and
 126 Kaufmann [9], the number of samples required scales with $\log(1/\delta)$, and so we aim to minimize

$$\limsup_{\delta \rightarrow 0} \frac{\mathbb{E}[\tau]}{\log(1/\delta)}$$

127 among algorithms with probability of error no more than δ . In this setting, we study the performance
 128 of EI when combined with the stopping rule studied by Chernoff [5] and Garivier and Kaufmann [9].

129 3 Sampling Rules

130 In this section, we first introduce the expected improvement algorithm, and point out its weakness.
 131 Then a simple variant of the expected improvement algorithm is proposed. Both algorithms make
 132 calculations using function $f(x) = x\Phi(x) + \phi(x)$ where $\Phi(\cdot)$ and $\phi(\cdot)$ are the CDF and PDF of
 133 the standard normal distribution. One can show that as $x \rightarrow \infty$, $\log f(-x) \sim -x^2/2$, and so
 134 $f(-x) \approx e^{-x^2/2}$ for very large x . One can also show that f is an increasing function.

135 **Expected Improvement.** *Expected improvement* [14] is a simple improvement-based sampling
 136 rule. The EI algorithm favors the arm that offers the largest amount of improvement upon a target.
 137 The EI algorithm measures the arm $I_n = \arg \max_{i \in A} v_{n,i}$ where $v_{n,i}$ is the EI value of arm i at time
 138 n . Let $I_n^* = \arg \max_{i \in A} \mu_{n,i}$ denote the arm with largest posterior mean at time n . The EI value of
 139 arm i at time n is defined as

$$v_{n,i} \triangleq \mathbb{E}_{\theta \sim \Pi_n} \left[(\theta_i - \mu_{n,I_n^*})^+ \right].$$

140 where $x^+ = \max\{x, 0\}$. The above expectation can be computed analytically as follows,

$$v_{n,i} = (\mu_{n,i} - \mu_{n,I_n^*}) \Phi \left(\frac{\mu_{n,i} - \mu_{n,I_n^*}}{\sigma_{n,i}} \right) + \sigma_{n,i} \phi \left(\frac{\mu_{n,i} - \mu_{n,I_n^*}}{\sigma_{n,i}} \right) = \sigma_{n,i} f \left(\frac{\mu_{n,i} - \mu_{n,I_n^*}}{\sigma_{n,i}} \right).$$

141 The EI value $v_{n,i}$ measures the potential of arm i to improve upon the largest posterior mean μ_{n,I_n^*} at
 142 time n . Because f is an increasing function, $v_{n,i}$ is increasing in both the posterior mean $\mu_{n,i}$ and
 143 posterior standard deviation $\sigma_{n,i}$.

144 **Top-Two Expected Improvement.** The EI algorithm can have very poor performance for selecting
 145 the best arm. Once it finds a particular arm with reasonably high probability to be the best, it allocates
 146 nearly all future samples to this arm at the expense of measuring other arms. Recently Ryzhov [22]
 147 showed that EI only allocates $\mathcal{O}(\log n)$ samples to suboptimal arms asymptotically. This is a severe
 148 shortcoming, as it means n must be extremely large before the algorithm has enough samples from
 149 suboptimal arms to reach a confident conclusion.

150 To improve the EI algorithm, we build on the top-two sampling idea in Russo [21]. The idea is to
 151 identify in each period the two ‘‘most promising’’ arms based on current observations, and randomize
 152 to choose which to sample. A tuning parameter $\beta \in (0, 1)$ controls the probability assigned to the
 153 ‘‘top’’ arm. A naive top-two variant of EI would identify the two arms with largest EI value, and flip
 154 a β -weighted coin to decide which to measure. However, one can prove that this algorithm is not
 155 optimal for any choice of β . Instead, what we call the top-two expected improvement algorithm uses
 156 a novel modified EI criterion which more carefully accounts for the decision-maker’s uncertainty
 157 when deciding which arm to sample.

158 For $i, j \in A$, define $v_{n,i,j} \triangleq \mathbb{E}_{\theta \sim \Pi_n} [(\theta_i - \theta_j)^+]$. This measures the expected magnitude of
 159 improvement arm i offers over arm j , but unlike the typical EI criterion, this expectation integrates
 160 over the uncertain quality of *both arms*. This measure can be computed analytically as

$$v_{n,i,j} = \sqrt{\sigma_{n,i}^2 + \sigma_{n,j}^2} f \left(\frac{\mu_{n,i} - \mu_{n,j}}{\sqrt{\sigma_{n,i}^2 + \sigma_{n,j}^2}} \right).$$

161 TTEI depends on a tuning parameter $\beta > 0$, set to $1/2$ by default. With probability β , TTEI measures
 162 the arm $I_n^{(1)}$ by optimizing the EI criterion, and otherwise it measures an alternative $I_n^{(2)}$ that offers
 163 the largest expected improvement on the arm $I_n^{(1)}$. Formally, TTEI measures the arm

$$I_n = \begin{cases} I_n^{(1)} = \arg \max_{i \in A} v_{n,i}, & \text{with probability } \beta, \\ I_n^{(2)} = \arg \max_{i \in A} v_{n,i,I_n^{(1)}}, & \text{with probability } 1 - \beta. \end{cases}$$

164 Note that $v_{n,i,i} = 0$, which implies $I_n^{(2)} \neq I_n^{(1)}$.

165 We notice that TTEI with $\beta = 1$ is the standard EI algorithm. Comparing to the EI algorithm, TTEI
 166 with $\beta \in (0, 1)$ allocates much more measurement effort to suboptimal arms. We will see that TTEI
 167 allocates β proportion of samples to the best arm asymptotically, and it uses the remaining $1 - \beta$
 168 fraction of samples for gathering evidence against each suboptimal arm.

169 4 Convergence to Asymptotically Optimal Proportions

170 For all $i \in A$ and $n \in \mathbb{N}$, we define $T_{n,i} \triangleq \sum_{\ell=1}^{n-1} \mathbf{1}\{I_\ell = i\}$ to be the number of samples of arm
 171 i before time n . We will show that under TTEI with parameter β , $\lim_{n \rightarrow \infty} T_{n,1}/n = \beta$. That is,
 172 the algorithm asymptotically allocates β proportion of the samples to true best arm. Dropping for
 173 the moment questions regarding the impact of this tuning parameter, let us consider the optimal
 174 asymptotic proportion of effort to allocate to each of the $k - 1$ remaining arms. It is known that the
 175 optimal proportions are given by the unique vector $(w_2^\beta, \dots, w_k^\beta)$ satisfying, $\sum_{i=2}^k w_i^\beta = 1 - \beta$ and

$$\frac{(\mu_2 - \mu_1)^2}{1/w_2^\beta + 1/\beta} = \dots = \frac{(\mu_k - \mu_1)^2}{1/w_k^\beta + 1/\beta}. \quad (2)$$

176 We set $w_1^\beta = \beta$, so $w^\beta = (w_1^\beta, \dots, w_k^\beta)$ encodes the sampling proportions of each arm.

177 To understand the source of equation (2), imagine that over the first n periods each arm i is sampled
 178 exactly $w_i^\beta n$ times, and let $\hat{\mu}_{n,i} \sim N\left(\mu_i, \frac{\sigma_i^2}{w_i^\beta n}\right)$ denote the empirical mean of arm i . Then

$$\hat{\mu}_{n,1} - \hat{\mu}_{n,i} \sim N\left(\mu_1 - \mu_i, \tilde{\sigma}_i^2\right) \quad \text{where} \quad \tilde{\sigma}_i^2 = \frac{\sigma^2}{n/\beta + n/w_i^\beta}.$$

179 The probability $\hat{\mu}_{n,1} - \hat{\mu}_{n,i} \leq 0$ —leading to an incorrect estimate of the arm with highest mean—is
 180 $\Phi((\mu_i - \mu_1)/\tilde{\sigma}_i)$ where Φ is the CDF of the standard normal distribution. Equation (2) is equivalent
 181 to requiring $(\mu_1 - \mu_i)/\tilde{\sigma}_i$ is equal for all arms i , so the probability of falsely declaring $\mu_i \geq \mu_1$
 182 is equal for all $i \neq 1$. In a sense, these sampling frequencies equalize the evidence against each
 183 suboptimal arm. These proportions appeared first in the machine learning literature in [21, 9], but
 184 appeared much earlier in the statistics literature in [13], and separately in the simulation optimization
 185 literature in [10]. As we will see in the next section, convergence to this allocation is a necessary
 186 condition for both notions of optimality considered in this paper.

187 Our main theoretical contribution is the following theorem, which establishes that under TTEI
 188 sampling proportions converge to the proportions w^β derived above. Therefore, while the sampling
 189 proportion of the optimal arm is controlled by the tuning parameter β , the remaining $1 - \beta$ fraction
 190 of measurement is optimally distributed among the remaining $k - 1$ arms. One of our results requires
 191 more than convergence to w^β with probability 1, but a sense in which the expected time until
 192 convergence is finite. To make this precise, we introduce a time after which for each arm, both its
 193 empirical mean and empirical proportion are accurate. Specifically, given $\beta \in (0, 1)$ and $\epsilon > 0$, we
 194 define

$$T_\beta^\epsilon \triangleq \inf \left\{ N \in \mathbb{N} : |\mu_{n,i} - \mu_i| \leq \epsilon \text{ and } |T_{n,i}/n - w_i^\beta| \leq \epsilon, \forall i \in A \text{ and } n \geq N \right\}. \quad (3)$$

195 If $T_{n,i}/n \rightarrow w_i^\beta$ with probability 1, then by the law of large numbers $\mathbb{P}(T_\beta^\epsilon < \infty) = 1$ for every
 196 $\epsilon > 0$. Such a result was established for other top-two sampling algorithms in [21]. To establish
 197 optimality in the “fixed confidence setting”, we need to prove in addition that $\mathbb{E}[T_\beta^\epsilon] < \infty$ for all
 198 $\epsilon > 0$, which requires substantial new technical innovations.

199 **Theorem 1.** *If TTEI is applied with parameter $\beta \in (0, 1)$, $\mathbb{E}[T_\beta^\epsilon] < \infty$ for any $\epsilon > 0$. Therefore,*

$$\lim_{n \rightarrow \infty} \frac{T_{n,i}}{n} = w_i^\beta \quad \forall i \in A.$$

200 4.1 Problem Complexity Measure

201 Given $\beta \in (0, 1)$, define the problem complexity measure

$$\Gamma_\beta^* \triangleq \frac{(\mu_2 - \mu_1)^2}{2\sigma^2 \left(1/w_2^\beta + 1/\beta\right)} = \dots = \frac{(\mu_k - \mu_1)^2}{2\sigma^2 \left(1/w_k^\beta + 1/\beta\right)},$$

202 which is a function of the true arm means and variances. This will be the exponent governing
 203 the rate of posterior convergence, and also characterizing the average number of samples in the

204 fixed confidence setting. The optimal exponent comes from maximizing over β . Let us define
 205 $\Gamma^* = \max_{\beta \in (0,1)} \Gamma_\beta^*$ and $\beta^* = \arg \max_{\beta \in (0,1)} \Gamma_\beta^*$ and set

$$w^* = w^{\beta^*} = \left(\beta^*, w_2^{\beta^*}, \dots, w_k^{\beta^*} \right).$$

206 Russo [21] has proved that for $\beta \in (0, 1)$, $\Gamma_\beta^* \geq \Gamma^* / \max \left\{ \frac{\beta^*}{\beta}, \frac{1-\beta^*}{1-\beta} \right\}$, and therefore $\Gamma_{1/2}^* \geq \Gamma^* / 2$.
 207 This demonstrates a surprising degree of robustness to β . In particular, Γ_β is close to Γ^* if β is
 208 adjusted to be close to β^* , and the choice of $\beta = 1/2$ always yields a 2-approximation to Γ^* .

209 5 Implied Optimality Results

210 This section establishes formal optimality guarantees for TTEI. Both results, in fact, hold for any
 211 algorithm satisfying the conclusions of Theorem 1, and is therefore one of broader interest.

212 5.1 Optimal Rate of Posterior Convergence

213 We first provide upper and lower bounds on the exponent governing the rate of posterior convergence.
 214 The same result has been proved in Russo [21] for bounded correlated priors. We use
 215 different proof techniques to prove the following result for uncorrelated Gaussian priors.

216 This theorem shows that no algorithm can attain a rate of posterior convergence faster than $e^{-\Gamma^* n}$
 217 and that this is attained by any algorithm that, like TTEI with optimal tuning parameter β^* , has
 218 asymptotic sampling ratios (w_1^*, \dots, w_k^*) . The second part implies TTEI with parameter β attains
 219 convergence rate $e^{-n\Gamma_\beta^*}$ and that it is optimal among sampling rules that allocation β -fraction of
 220 samples to the optimal arm. Recall that, without loss of generality, we have assumed arm 1 is the arm
 221 with true highest mean $\mu_1 = \max_{i \in A} \mu_i$. We will study the posterior mass $1 - \alpha_{n,1}$ assigned to the
 222 event that some other has the highest mean.

223 **Theorem 2** (Posterior Convergence - Sufficient Condition for Optimality). *The following properties*
 224 *hold with probability 1:*

225 1. Under any allocation rule satisfying $T_{n,i}/n \rightarrow w_i^*$ for each $i \in A$,

$$\lim_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,1}) = \Gamma^*.$$

226 Under any sampling rule,

$$\limsup_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,1}) \leq \Gamma^*.$$

227 2. For $\beta \in (0, 1)$, under any allocation rule satisfying $T_{n,i}/n \rightarrow w_i^\beta$ for each $i \in A$,

$$\lim_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,1}) = \Gamma_\beta^*.$$

228 Under any sampling rule satisfying $T_{n,1}/n \rightarrow \beta$,

$$\limsup_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,1}) \leq \Gamma_\beta^*.$$

229 This result reveals that when the tuning parameter β is set optimally to β^* , TTEI attains the optimal
 230 rate of posterior convergence. Since $\Gamma_{1/2}^* \geq \Gamma^* / 2$, when β set to the default value $1/2$, the exponent
 231 governing the convergence rate of TTEI is at least half of the optimal one.

232 5.2 Optimal Average Sample Size

233 **Chernoff's Stopping Rule.** In the fixed confidence setting, besides an efficient sampling rule, a
 234 player also needs to design an intelligent stopping rule. This section introduces a stopping rule
 235 proposed by Chernoff [5] and studied recently by Garivier and Kaufmann [9]. This stopping rule
 236 makes use of the Generalized Likelihood Ratio statistic, which depends on the current maximum

237 likelihood estimates of all unknown means. For each arm $i \in A$, the maximum likelihood estimate of
 238 its unknown mean μ_i at time n is its empirical mean $\hat{\mu}_{n,i} = T_{n,i}^{-1} \sum_{\ell=1}^{n-1} \mathbf{1}\{I_\ell = i\} Y_{\ell, I_\ell}$. If $T_{n,i} = 0$,
 239 we set $\hat{\mu}_{n,i} = 0$. For arms $i, j \in A$, if $\hat{\mu}_{n,i} \geq \hat{\mu}_{n,j}$, the Generalized Likelihood Ratio statistic $Z_{n,i,j}$
 240 has the following explicit expression for Gaussian noise distributions:

$$Z_{n,i,j} \triangleq T_{n,i} d(\hat{\mu}_{n,i}, \hat{\mu}_{n,i,j}) + T_{n,j} d(\hat{\mu}_{n,j}, \hat{\mu}_{n,i,j})$$

241 where $d(x, y) \triangleq (x - y)^2 / (2\sigma^2)$ is the KL-divergence between two normal distributions $N(x, \sigma^2)$
 242 and $N(y, \sigma^2)$, and $\hat{\mu}_{n,i,j}$ is a weighted average of the empirical means of arms i, j defined as

$$\hat{\mu}_{n,i,j} \triangleq \frac{T_{n,i}}{T_{n,i} + T_{n,j}} \hat{\mu}_{n,i} + \frac{T_{n,j}}{T_{n,i} + T_{n,j}} \hat{\mu}_{n,j}.$$

243 On the other hand, if $\hat{\mu}_{n,i} < \hat{\mu}_{n,j}$, then $Z_{n,j,i}$ is well-defined as above, and $Z_{n,i,j} = -Z_{n,j,i} \leq 0$ (if
 244 $T_{n,i} = T_{n,j} = 0$, we let $Z_{n,i,j} = Z_{n,j,i} = 0$). Given a target confidence $\delta \in (0, 1)$, to ensure that
 245 one arm is better than the others with probability at least $1 - \delta$, we use the stopping time

$$\tau_\delta \triangleq \inf \left\{ n \in \mathbb{N} : Z_n \triangleq \max_{i \in A} \min_{j \in A \setminus \{i\}} Z_{n,i,j} > \gamma_{n,\delta} \right\}$$

246 where $\gamma_{n,\delta} > 0$ is an appropriate threshold. By definition, we know that $\min_{j \in A \setminus \{i\}} Z_{n,i,j}$ is
 247 nonnegative if and only if $\hat{\mu}_{n,i} \geq \hat{\mu}_{n,j}$ for all $j \in A \setminus \{i\}$. Hence, whenever $\hat{I}_n^* \triangleq \arg \max_{i \in A} \hat{\mu}_{n,i}$
 248 is unique, $Z_n = \min_{j \in A \setminus \{\hat{I}_n^*\}} Z_{n,\hat{I}_n^*,j}$.

249 Next we introduce the exploration rate for normal bandit models that can ensure to identify the best
 250 arm with probability at least $1 - \delta$. We use the following result given in Garivier and Kaufmann [9].

251 **Proposition 1** (Garivier and Kaufmann [9] Proposition 12). *Let $\delta \in (0, 1)$ and $\alpha > 1$. For any*
 252 *normal bandit model, there exists a constant $C = C(\alpha, k)$ such that under any possible sampling*
 253 *rule, using the Chernoff's stopping rule with the threshold $\gamma_{n,\delta}^\alpha = \log(Cn^\alpha / \delta)$ guarantees*

$$\mathbb{P} \left(\tau_\delta < \infty, \arg \max_{i \in A} \hat{\mu}_{\tau_\delta, i} \neq 1 \right) \leq \delta.$$

254 **Sample Complexity.** Garivier and Kaufmann [9] recently provided a general lower bound on the
 255 number of samples required in the fixed confidence setting. In particular, they show that for any
 256 normal bandit model, under any sampling rule and stopping time τ_δ that guarantees a probability of
 257 error less than δ ,

$$\liminf_{\delta \rightarrow 0} \frac{\mathbb{E}[\tau_\delta]}{\log(1/\delta)} \geq \frac{1}{\Gamma^*}.$$

258 Recall that T_β^ϵ , defined in (3), is the first time after which the empirical means and empirical
 259 proportions are within ϵ of their asymptotic limits. The next result provides a condition in terms of
 260 T_β^ϵ that is sufficient to guarantees optimality in the fixed confidence setting.

261 **Theorem 3** (Fixed Confidence - Sufficient Condition for Optimality). *Let $\beta \in (0, 1)$. Consider any*
 262 *sampling rule which, if applied with no stopping rule, satisfies $\mathbb{E}[T_\beta^\epsilon] < \infty$ for all $\epsilon > 0$. Fix any*
 263 *$\alpha > 1$. Then if this sampling rule is applied with Chernoff's stopping rule with the threshold $\gamma_{n,\delta}^\alpha$,*
 264 *we have*

$$\limsup_{\delta \rightarrow 0} \frac{\mathbb{E}[\tau_\delta]}{\log(1/\delta)} \leq \frac{\alpha}{\Gamma_\beta^*}.$$

265 Since α can be chosen to be arbitrarily close to 1, when $\beta = \beta^*$ the general lower bound on sample
 266 complexity of $1/\Gamma^*$ is essentially matched. In addition, when β is set to the default value $1/2$ and α
 267 is taken to be arbitrarily close to 1, the sample complexity of TTEI combined with the Chernoff's
 268 stopping rule is at most twice the optimal sample complexity since $1/\Gamma_{1/2}^* \leq 2/\Gamma^*$.

269 6 Numerical Experiments

270 To test the empirical performances of TTEI, we conduct several numerical experiments. The first
 271 experiment compares the performance of TTEI with $\beta = 1/2$ and EI. The second experiment

272 compares the performances of different versions of TTEI, top-two Thompson sampling (TTTS) [21],
 273 knowledge gradient (KG) [7] and oracle algorithms that know the optimal proportions *a priori*. Each
 274 algorithm plays arm $i = 1, \dots, k$ exactly once at the beginning, and then prescribe a prior $N(Y_{i,i}, \sigma^2)$
 275 for unknown arm-mean μ_i where $Y_{i,i}$ is the observation from $N(\mu_i, \sigma^2)$. In both experiments, we fix
 276 the common known variance $\sigma^2 = 1$ and the number of arms $k = 5$. We consider three instances
 277 $[\mu_1, \dots, \mu_5] = [5, 4, 1, 1, 1], [5, 4, 3, 2, 1]$ and $[2, 0.8, 0.6, 0.4, 0.2]$. The optimal parameter β^* equals
 278 0.48, 0.45 and 0.35, respectively.

279 Recall that $\alpha_{n,i}$, defined in (1), denotes the posterior probability that arm i is optimal. Table 1 shows
 280 the average number of measurements required for the largest posterior probability being the best to
 281 reach a given confidence level c , i.e., $\max_i \alpha_{n,i} \geq c$. The results in Table 1 are averaged over 100
 trials. We see that TTEI with $\beta = 1/2$ outperforms standard EI by an order of magnitude.

Table 1: Average number of measurements required to reach the confidence level $c = 0.95$

| | TTEI-1/2 | EI |
|-----------------------|----------|---------|
| $[5, 4, 1, 1, 1]$ | 14.60 | 238.50 |
| $[5, 4, 3, 2, 1]$ | 16.72 | 384.73 |
| $[2, .8, .6, .4, .2]$ | 24.39 | 1525.42 |

282

283 The second experiment compares the performance of different versions of TTEI, TTTS, KG, random
 284 sampling oracle (RSO) and tracking oracle (TO). The random sampling oracle draws a random arm in
 285 each round from the distribution w^* encoding the asymptotically optimal proportions. The tracking
 286 oracle tracks the optimal proportions at each round. Specifically, the tracking oracle samples the arm
 287 with the largest ratio its optimal and empirical proportions. Two tracking algorithms proposed by
 288 Garivier and Kaufmann [9] are similar to this tracking oracle. TTEI with adaptive β (aTTEI) works
 289 as follows: it starts with $\beta = 1/2$ and updates $\beta = \hat{\beta}^*$ every 10 rounds where $\hat{\beta}^*$ is the maximizer of
 290 equation (2) based on plug-in estimators for the unknown arm-means. Table 2 shows the average
 291 number of measurements required for the largest posterior probability being the best to reach the
 292 confidence level $c = 0.9999$. The results in Table 2 are averaged over 200 trials. We see that the
 293 performances of TTEI with adaptive β and TTEI with β^* are better than the performances of all other
 294 algorithms. We note that TTEI with adaptive β substantially outperforms the tracking oracle.

Table 2: Average number of measurements required to reach the confidence level $c = 0.9999$

| | TTEI-1/2 | aTTEI | TTEI- β^* | TTTS- β^* | RSO | TO | KG |
|-----------------------|----------|--------------|-----------------|-----------------|--------|-------|-------|
| $[5, 4, 1, 1, 1]$ | 61.97 | 61.98 | 61.59 | 62.86 | 97.04 | 77.76 | 75.55 |
| $[5, 4, 3, 2, 1]$ | 66.56 | 65.54 | 65.55 | 66.53 | 103.43 | 88.02 | 81.49 |
| $[2, .8, .6, .4, .2]$ | 76.21 | 72.94 | 71.62 | 73.02 | 101.97 | 96.90 | 86.98 |

295 7 Conclusion and Extensions to Correlated Arms

296 We conclude by noting that while this paper thoroughly studies TTEI in the case of uncorrelated
 297 priors, we believe the algorithm is also ideally suited to problems with complex correlated priors
 298 and large sets of arms. In fact, the modified information measure $v_{n,i,j}$ was designed with an eye
 299 toward dealing with correlation in a sophisticated way. In the case of a correlated normal distribution
 300 $N(\mu, \Sigma)$, one has

$$v_{n,i,j} = \mathbb{E}_{\theta \sim N(\mu, \Sigma)}[(\theta_i - \theta_j)^+] = \sqrt{\Sigma_{ii} + \Sigma_{jj} - 2\Sigma_{ij}} f\left(\frac{\mu_{n,i} - \mu_{n,j}}{\sqrt{\Sigma_{ii} + \Sigma_{jj} - 2\Sigma_{ij}}}\right).$$

301 This closed form accommodates efficient computation. Here the term $\Sigma_{i,j}$ accounts for the correlation
 302 or similarity between arms i and j . Therefore $v_{n,i,I_n^{(1)}}$ is large for arms i that offer large potential
 303 improvement over $I_n^{(1)}$, i.e. those that (1) have large posterior mean, (2) have large posterior variance,
 304 and (3) are not highly correlated with arm $I_n^{(1)}$. As $I_n^{(1)}$ concentrates near the estimated optimum, we
 305 expect the third factor will force the algorithm to experiment in promising regions of the domain that
 306 are “far” away from the current-estimated optimum, and are under-explored under standard EI.

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386 A Outline

387 The appendix is organized as follows.

- 388 1. Section B introduces some further notations required in the theoretical analysis.
- 389 2. Section C is the proof of Theorem 2, a sufficient condition in terms of optimal proportions
390 $(w_1^\beta, \dots, w_k^\beta)$ to guarantee the optimal rate of posterior convergence.
- 391 3. Section D is the proof of Theorem 3, a sufficient condition in terms of T_β^ϵ under which the
392 optimality in the fixed confidence setting is achieved.
- 393 4. Section E provides several basic results which is used in the theoretical analysis of TTEI.
- 394 5. Section F proves that TTEI satisfies the sufficient conditions for two notions of optimality,
395 which immediately establishes Theorems 1.

396 B Notation

397 For notational convenience, we assume that sampling rules begin with an improper prior for each arm
398 $i \in A$ with $\mu_{1,i} = 0$ and $\sigma_{1,i} = \infty$. Consequently, if $T_{n,i} = \sum_{\ell=1}^{n-1} \mathbf{1}\{I_\ell = i\} = 0$, $\mu_{n,i} = \mu_{1,i} = 0$
399 and $\sigma_{n,i} = \sigma_{1,i} = \infty$, and if $T_{n,i} > 0$,

$$\mu_{n,i} = \frac{1}{T_{n,i}} \sum_{\ell=1}^{n-1} \mathbf{1}\{I_\ell = i\} Y_{\ell, I_\ell} \quad \text{and} \quad \sigma_{n,i}^2 = \frac{\sigma^2}{T_{n,i}},$$

400 so the posterior parameters are identical to the frequentist sample mean and variance under the
401 observations collected so far.

402 We introduce some further notations. We define

$$\Delta_{\min} \triangleq \min_{i \neq j} |\mu_i - \mu_j| \quad \text{and} \quad \Delta_{\max} \triangleq \max_{i, j \in A} (\mu_i - \mu_j).$$

403 Since the arm means are unique, we have $\Delta_{\min}, \Delta_{\max} > 0$. In addition, we define

$$\beta_{\min} \triangleq \min\{\beta, 1 - \beta\} \quad \text{and} \quad \beta_{\max} \triangleq \max\{\beta, 1 - \beta\}.$$

404 Note that for $\beta \in (0, 1)$, $\beta_{\min} > 0$.

405 We introduce the filtration $(\mathcal{F}_n : n = 1, 2, \dots)$ where

$$\mathcal{F}_n = \Sigma(I_1, Y_{1, I_1}, \dots, I_n, Y_{n, I_n})$$

406 is the sigma algebra generated by observations up to time n . For all $i \in A$ and $n \in \mathbb{N}$, define

$$\psi_{n,i} \triangleq \mathbb{P}(I_n = i | \mathcal{F}_{n-1}) \quad \text{and} \quad \Psi_{n,i} \triangleq \sum_{\ell=1}^{n-1} \psi_{\ell, i}.$$

407 Note that for all $i \in A$, $T_{1,i} = \Psi_{1,i} = 0$. Both $T_{n,i}$ and $\Psi_{n,i}$ measure the effort allocated to arm i up
408 to period n .

409 Finally, rather than use the notation $v_{n,i}$ and $v_{n,i,j}$ introduced in Section 3 for the expected-
410 improvement measures it is more convenient to work with the notation defined here. Set

$$v_{n,i}^{(1)} \equiv v_{n,i} \quad \forall i \in A$$

411 to be the expected improvement used in the identifying the first among in the top-two, and

$$v_{n,i}^{(2)} \equiv v_{n,i, I_n^{(1)}} \quad \forall i \in A$$

412 to be the second expected improvement measure where $I_n^{(1)}$ is the arm optimizing the first expected
413 improvement measure.

414 **C Proof of Theorem 2**

415 To prove Theorem 2, we first need to introduce the so-called Gaussian tail inequality.

416 **Lemma 1.** *Let $X \sim N(\mu, \sigma^2)$ and $c \geq 0$, then we have*

$$\frac{1}{\sqrt{2\pi}} e^{-(\sigma+c)^2/(2\sigma^2)} \leq \mathbb{P}(X \geq \mu + c) \leq \frac{1}{2} e^{-c^2/(2\sigma^2)}.$$

417 *Proof.* We first prove the upper bound.

$$\begin{aligned} \mathbb{P}(X \geq \mu + c) &= \int_{\mu+c}^{\infty} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(x-\mu)^2/(2\sigma^2)} dx \\ &= \int_0^{\infty} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(x+c)^2/(2\sigma^2)} dx \\ &\leq \int_0^{\infty} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(x^2+c^2)/(2\sigma^2)} dx \\ &= e^{-c^2/(2\sigma^2)} \int_0^{\infty} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-x^2/(2\sigma^2)} dx \\ &= \frac{1}{2} e^{-c^2/(2\sigma^2)}. \end{aligned}$$

418 Next we prove the lower bound.

$$\begin{aligned} \mathbb{P}(X \geq \mu + c) &= \int_{\mu+c}^{\infty} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(x-\mu)^2/(2\sigma^2)} dx \\ &= \int_0^{\infty} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(x+c)^2/(2\sigma^2)} dx \\ &\geq \int_0^{\sigma} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(x+c)^2/(2\sigma^2)} dx \\ &\geq \int_0^{\sigma} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(\sigma+c)^2/(2\sigma^2)} dx \\ &= \frac{1}{\sqrt{2\pi}} e^{-(\sigma+c)^2/(2\sigma^2)}. \end{aligned}$$

419

□

420 **Proof of Theorem 2.** We let $\mathcal{I} = \{i \in A : \lim_{n \rightarrow \infty} T_{n,i} = \infty\}$ and $\bar{\mathcal{I}} = A \setminus \mathcal{I}$. Note that $\bar{\mathcal{I}}$
421 contains arms that are only sampled finite times. First, suppose that $\bar{\mathcal{I}}$ is nonempty. For each $i \in A$,
422 we define

$$\mu_{\infty,i} \triangleq \lim_{n \rightarrow \infty} \mu_{n,i} \quad \text{and} \quad \sigma_{\infty,i}^2 \triangleq \lim_{n \rightarrow \infty} \sigma_{n,i}^2.$$

423 Recall that for each $i \in A$, an improper prior with $\mu_{1,i} = 0$ and $\sigma_{1,i} = \infty$ is prescribed. Then if
424 $T_{n,i} = \sum_{\ell=1}^{n-1} \mathbf{1}\{I_{\ell} = i\} = 0$, $\mu_{n,i} = \mu_{1,i} = 0$ and $\sigma_{n,i} = \sigma_{1,i} = \infty$, and if $T_{n,i} > 0$.

$$\mu_{n,i} = \frac{1}{T_{n,i}} \sum_{\ell=1}^{n-1} \mathbf{1}\{I_{\ell} = i\} Y_{\ell, I_{\ell}} \quad \text{and} \quad \sigma_{n,i}^2 = \frac{\sigma^2}{T_{n,i}},$$

425 Hence, for $i \in \mathcal{I}$, $\mu_{\infty,i} = \mu_i$ and $\sigma_{\infty,i}^2 = 0$, while for $i \in \bar{\mathcal{I}}$, $\sigma_{\infty,i}^2 > 0$. We let

$$\Pi_{\infty} = N(\mu_{\infty,1}, \sigma_{\infty,1}^2) \otimes N(\mu_{\infty,2}, \sigma_{\infty,2}^2) \otimes \cdots \otimes N(\mu_{\infty,k}, \sigma_{\infty,k}^2),$$

426 and for each $i \in A$, we define

$$\alpha_{\infty,i} \triangleq \mathbb{P}_{\theta \sim \Pi_{\infty}} \left(\theta_i > \max_{j \neq i} \theta_j \right).$$

427 For $i \in \bar{\mathcal{I}}$ is nonempty, we have $\alpha_{\infty,i} \in (0, 1)$ since $\sigma_{\infty,i}^2 > 0$. This implies $\alpha_{\infty,1} < 1$ and so

$$\lim_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,1}) = \lim_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{\infty,1}) = 0.$$

428 Now suppose $\bar{\mathcal{I}}$ is empty. By definition, $\alpha_{n,1} = \mathbb{P}_{\theta \sim \Pi_n}(\theta_1 > \max_{i \neq 1} \theta_i)$, so $1 - \alpha_{n,1} =$
429 $\mathbb{P}_{\theta \sim \Pi_n}(\cup_{i \neq 1}(\theta_i \geq \theta_1))$, and then we have

$$\max_{i \neq 1} \mathbb{P}_{\theta \sim \Pi_n}(\theta_i \geq \theta_1) \leq 1 - \alpha_{n,1} \leq \sum_{i \neq 1} \mathbb{P}_{\theta \sim \Pi_n}(\theta_i \geq \theta_1) \leq (k-1) \max_{i \neq 1} \mathbb{P}_{\theta \sim \Pi_n}(\theta_i \geq \theta_1) \quad (4)$$

430 where the second inequality uses the union bound.

431 To simplify the presentation, we need to introduce the following asymptotic notation. We say two
432 real-valued sequences $\{a_n\}$ and $\{b_n\}$ are *logarithmically equivalent* if $\lim_{n \rightarrow \infty} 1/n \log(a_n/b_n) = 0$.

433 We denote this by $a_n \doteq b_n$. Using equation 4, we conclude

$$1 - \alpha_{n,1} \doteq \max_{i \neq 1} \mathbb{P}_{\theta \sim \Pi_n}(\theta_i \geq \theta_1).$$

434 Next we want to show that for $i \neq 1$, $\mathbb{P}_{\theta \sim \Pi_n}(\theta_i \geq \theta_1) \doteq \exp\left(\frac{-(\mu_{n,i} - \mu_{n,1})^2}{2\sigma^2(1/T_{n,i} + 1/T_{n,1})}\right)$. Note that at
435 time n , $\theta_i - \theta_1 \sim N(\mu_{n,i} - \mu_{n,1}, \sigma_{n,i}^2 + \sigma_{n,1}^2)$ and $\sigma_{n,i}^2 + \sigma_{n,1}^2 = \sigma^2(1/T_{n,i} + 1/T_{n,1})$. Since every
436 arm is sampled infinite times, when n is large, $\mu_{n,1} \geq \mu_{n,i}$, and then using Lemma 1, we have

$$\frac{1}{\sqrt{2\pi}} \exp\left(\frac{-\left(\sqrt{\sigma_{n,i}^2 + \sigma_{n,1}^2} + \mu_{n,1} - \mu_{n,i}\right)^2}{2(\sigma_{n,i}^2 + \sigma_{n,1}^2)}\right) \leq \mathbb{P}_{\theta \sim \Pi_n}(\theta_i - \theta_1 \geq 0) \leq \frac{1}{2} \exp\left(\frac{-(\mu_{n,1} - \mu_{n,i})^2}{2(\sigma_{n,i}^2 + \sigma_{n,1}^2)}\right),$$

437 which implies

$$\frac{1}{n} \log\left(\frac{1}{\sqrt{2\pi}}\right) - \frac{1}{2n} - \frac{\mu_{n,1} - \mu_{n,i}}{n\sqrt{\sigma_{n,i}^2 + \sigma_{n,1}^2}} \leq \frac{1}{n} \log\left(\frac{\mathbb{P}_{\theta \sim \Pi_n}(\theta_i \geq \theta_1)}{\exp\left(\frac{-(\mu_{n,1} - \mu_{n,i})^2}{2(\sigma_{n,i}^2 + \sigma_{n,1}^2)}\right)}\right) \leq \frac{1}{n} \log\left(\frac{1}{2}\right).$$

438 Note that when $\mu_{n,1} \geq \mu_{n,i}$,

$$0 \leq \frac{\mu_{n,1} - \mu_{n,i}}{n\sqrt{\sigma_{n,i}^2 + \sigma_{n,1}^2}} = \frac{\mu_{n,1} - \mu_{n,i}}{\sigma\sqrt{n(n/T_{n,i} + n/T_{n,1})}} \leq \frac{\mu_{n,1} - \mu_{n,i}}{\sigma\sqrt{2n}}$$

439 where the last equality uses $T_{n,i}, T_{n,1} < n$. Using the squeeze theorem, we have

$$\lim_{n \rightarrow \infty} \frac{\mu_{n,1} - \mu_{n,i}}{n\sqrt{\sigma_{n,i}^2 + \sigma_{n,1}^2}} = 0,$$

440 and

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log\left(\frac{\mathbb{P}_{\theta \sim \Pi_n}(\theta_i \geq \theta_1)}{\exp\left(\frac{-(\mu_{n,1} - \mu_{n,i})^2}{2(\sigma_{n,i}^2 + \sigma_{n,1}^2)}\right)}\right) = 0.$$

441 Hence, $\mathbb{P}_{\theta \sim \Pi_n}(\theta_i \geq \theta_1) \doteq \exp\left(\frac{-(\mu_{n,i} - \mu_{n,1})^2}{2\sigma^2(1/T_{n,i} + 1/T_{n,1})}\right)$. Then we have

$$\begin{aligned} 1 - \alpha_{n,i} &\doteq \max_{i \neq 1} \mathbb{P}_{\theta \sim \Pi_n}(\theta_i \geq \theta_1) \\ &\doteq \max_{i \neq 1} \left\{ \exp\left(\frac{-(\mu_{n,i} - \mu_{n,1})^2}{2\sigma^2(1/T_{n,i} + 1/T_{n,1})}\right) \right\} \\ &\doteq \exp\left(-n \min_{i \neq 1} \left\{ \frac{(\mu_{n,i} - \mu_{n,1})^2}{2\sigma^2(n/T_{n,i} + n/T_{n,1})} \right\}\right) \end{aligned}$$

442 where the second equality uses the property that if $a_{n,i} \doteq b_{n,i}$ for each $i = 1, \dots, c$ where c a positive
443 integer, then $\max_{i \in \{1, \dots, c\}} a_{n,i} \doteq \max_{i \in \{1, \dots, c\}} b_{n,i}$.

444 Let $W \triangleq \left\{ w = (w_1, \dots, w_k) : \sum_{i=1}^k w_i = 1 \text{ and } w_i \geq 0, \forall i \in A \right\}$ denote the set of possible pro-
 445 portions on k arms. Russo [21] showed that

$$\Gamma^* = \max_{w \in W} \min_{i \neq 1} \frac{(\mu_i - \mu_1)^2}{2\sigma^2(1/w_i + 1/w_1)},$$

446 and given $\beta \in (0, 1)$,

$$\Gamma_\beta^* = \max_{w \in W: w_1 = \beta} \min_{i \neq 1} \frac{(\mu_i - \mu_1)^2}{2\sigma^2(1/w_i + 1/w_1)}.$$

447 Under any sampling rule,

$$\begin{aligned} 1 - \alpha_{n,i} &\doteq \exp \left(-n \min_{i \neq 1} \left\{ \frac{(\mu_{n,i} - \mu_{n,1})^2}{2\sigma^2(n/T_{n,i} + n/T_{n,1})} \right\} \right) \\ &\geq \exp \left(-n \max_{w \in W} \min_{i \neq 1} \left\{ \frac{(\mu_{n,i} - \mu_{n,1})^2}{2\sigma^2(1/w_i + 1/w_1)} \right\} \right) \end{aligned}$$

448 Since every arm is sampled infinite times, as $n \rightarrow \infty$, $\mu_{n,i} \rightarrow \mu_i$ and $\mu_{n,1} \rightarrow \mu_1$, and thus

$$\limsup_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,i}) \leq \Gamma^*.$$

449 If $T_{n,i}/n \rightarrow w_i^*$ for each $i \in A$, then for each $i \neq 1$, we have

$$\lim_{n \rightarrow \infty} \frac{(\mu_{n,i} - \mu_{n,1})^2}{2\sigma^2(n/T_{n,i} + n/T_{n,1})} = \frac{(\mu_i - \mu_1)^2}{2\sigma^2(1/w_i^* + 1/\beta)} = \Gamma^*,$$

450 and thus

$$1 - \alpha_{n,i} \doteq \exp \left(-n \min_{i \neq 1} \left\{ \frac{(\mu_{n,i} - \mu_{n,1})^2}{2\sigma^2(n/T_{n,i} + n/T_{n,1})} \right\} \right) \doteq \exp(-n\Gamma^*),$$

451 which implies

$$\lim_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,i}) = \Gamma^*.$$

452 Similarly, for $\beta \in (0, 1)$, under any sampling rule satisfying $T_{n,1}/n \rightarrow \beta$, we have

$$\limsup_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,i}) \leq \Gamma_\beta^*,$$

453 and under any sampling rule satisfying $T_{n,i}/n \rightarrow w_i^\beta$ for each $i \in A$,

$$\lim_{n \rightarrow \infty} -\frac{1}{n} \log(1 - \alpha_{n,i}) = \Gamma_\beta^*.$$

454 D Proof of Theorem 3

455 Let $\beta \in (0, 1)$. Recall that TTEI begins with an improper prior for each arm $i \in A$ with $\mu_{1,i} = 0$ and
 456 $\sigma_{1,i} = \infty$, so for any $i \in A$ and $n \in \mathbb{N}$, $\mu_{n,i} = \hat{\mu}_{n,i}$, i.e., the posterior mean equals the empirical
 457 mean, and thus $I_n^* = \arg \max_{i \in A} \mu_{n,i}$ is identical to $\hat{I}_n^* = \arg \max_{i \in A} \hat{\mu}_{n,i}$. We can rewrite Z_n
 458 used in the Chernoff's stopping rule as follows,

$$Z_n = \min_{j \in A \setminus \{I_n^*\}} Z_{n, I_n^*, j}$$

459 where the Generalized Likelihood Ratio statistic is

$$Z_{n, I_n^*, j} = T_{n, I_n^*} d(\mu_{n, I_n^*}, \mu_{n, I_n^*, j}) + T_{n, j} d(\mu_{n, j}, \mu_{n, I_n^*, j})$$

460 where

$$\mu_{n, I_n^*, j} = \frac{T_{n, I_n^*}}{T_{n, I_n^*} + T_{n, j}} \mu_{n, I_n^*} + \frac{T_{n, j}}{T_{n, I_n^*} + T_{n, j}} \mu_{n, j}.$$

461 Note that $\Delta_{\min} = \min_{i \neq j} |\mu_i - \mu_j| > 0$. Then by definition of $T_\beta^{\Delta_{\min}/4}$, for all $i \in A$ and
 462 $n \geq T_\beta^{\Delta_{\min}/4}$, $|\mu_{n,i} - \mu_i| \leq \Delta_{\min}/4$, which implies $\mu_{n,1} > \dots > \mu_{n,k}$, and thus $I_n^* = 1$. Using
 463 $d(x, y) = (x - y)^2 / (2\sigma^2)$, for $n \geq T_\beta^{\Delta_{\min}/4}$, we have

$$\frac{Z_n}{n} = \min_{i \in A \setminus \{1\}} \frac{(\mu_{n,i} - \mu_{n,1})^2}{2\sigma^2(n/T_{n,i} + n/T_{n,1})}.$$

464 Note that

$$\Gamma_\beta^* = \frac{(\mu_2 - \mu_1)^2}{2\sigma^2(1/w_2^\beta + 1/\beta)} = \dots = \frac{(\mu_k - \mu_1)^2}{2\sigma^2(1/w_k^\beta + 1/\beta)}$$

465 and when $\beta \in (0, 1)$, $\Gamma_\beta^* > 0$. Given $\epsilon > 0$, there exists $\epsilon' \in (0, \Delta_{\min}/4]$ such that for all
 466 $n \geq N^\epsilon \triangleq T_\beta^{\epsilon'}$, $|\mu_{n,i} - \mu_i| \leq \epsilon'$ and $|T_{n,i}/n - w_i^\beta| \leq \epsilon', \forall i \in A$ can imply $Z_n/n \geq \Gamma_\beta^* - \epsilon$. We
 467 have $\mathbb{E}[N^\epsilon] = \mathbb{E}[T_\beta^{\epsilon'}] < \infty$.

468 Let $\delta \in (0, 1)$ and $\alpha > 0$. By Proposition 1, the stopping time $\tau_\delta =$
 469 $\inf\{n \in \mathbb{N} : Z_n > \log(Cn^\alpha/\delta)\}$ can ensure $\mathbb{P}(\tau_\delta < \infty, \arg \max_{i \in A} \mu_{\tau_\delta, i} \neq 1) \leq \delta$.

470 For $\epsilon \in (0, \Gamma_\beta^*/(1 + \alpha))$, when $n \geq N^\epsilon$, $Z_n \geq (\Gamma_\beta^* - \epsilon)n > 0$. Let $M^\epsilon \triangleq \lceil \max\{N^\epsilon, 1/\epsilon^2\} \rceil$
 471 where the ceil function $\lceil x \rceil$ is the least integer greater than or equal to x . Now let us consider the
 472 following two cases.

473 1. $\exists r \in [1, M^\epsilon]$ such that $Z_r > \log(Cr^\alpha/\delta)$

474 This case implies $\tau_\delta \leq M^\epsilon$.

475 2. $\forall r \in [1, M^\epsilon]$, $Z_r \leq \log(Cr^\alpha/\delta)$

476 This case implies $\tau_\delta \geq M^\epsilon + 1$. Note that $M^\epsilon = \lceil \max\{N^\epsilon, 1/\epsilon^2\} \rceil \geq N^\epsilon$, so for
 477 $n \geq M^\epsilon$, $Z_n \geq (\Gamma_\beta^* - \epsilon)n$. Let x^ϵ be the solution of $(\Gamma_\beta^* - \epsilon)x = \log(Cx^\alpha/\delta)$. Since
 478 $(\Gamma_\beta^* - \epsilon)M^\epsilon \leq Z_{M^\epsilon} \leq \log(C(M^\epsilon)^\alpha/\delta)$, we have $x^\epsilon \geq M^\epsilon$, which implies $x^\epsilon \geq 1/\epsilon^2$, and
 479 then $\log(x^\epsilon) \leq (x^\epsilon)^{1/2} \leq \epsilon x^\epsilon$. Hence, $(\Gamma_\beta^* - \epsilon)x^\epsilon = \log(C(x^\epsilon)^\alpha/\delta) \leq \log(C) + \alpha \epsilon x^\epsilon +$
 480 $\log(1/\delta)$, which implies

$$x^\epsilon \leq \frac{\log(C) + \log(1/\delta)}{\Gamma_\beta^* - (1 + \alpha)\epsilon}.$$

481 Let $L_\delta^\epsilon \triangleq \inf\{n \geq M^\epsilon : (\Gamma_\beta^* - \epsilon)n > \log(Cn^\alpha/\delta)\}$. Since $(\Gamma_\beta^* - \epsilon)x^\epsilon =$
 482 $\log(C(x^\epsilon)^\alpha/\delta)$, we have

$$L_\delta^\epsilon \leq \lceil x^\epsilon \rceil + 1 \leq \left\lceil \frac{\log(C) + \log(1/\delta)}{\Gamma_\beta^* - (1 + \alpha)\epsilon} \right\rceil + 1 < \frac{\log(C) + \log(1/\delta)}{\Gamma_\beta^* - (1 + \alpha)\epsilon} + 2.$$

483 We notice that $Z_{L_\delta^\epsilon} \geq (\tau_\beta^* - \epsilon)L_\delta^\epsilon > \log(C(L_\delta^\epsilon)^\alpha/\delta)$, so we have $\tau_\delta \leq L_\delta^\epsilon$.

484 Combining the above two cases, we have $\tau_\delta \leq M^\epsilon + L_\delta^\epsilon$, and thus $\mathbb{E}[\tau_\delta] \leq \mathbb{E}[M^\epsilon] + \mathbb{E}[L_\delta^\epsilon]$. Note
 485 that $M^\epsilon = \lceil \max\{N^\epsilon, 1/\epsilon^2\} \rceil$ and $\mathbb{E}[N^\epsilon] < \infty$ imply $\mathbb{E}[M^\epsilon] < \infty$.

486 Now we fix $\tilde{\epsilon} = (\alpha - 1)\Gamma_\beta^*/[\alpha(1 + \alpha)] \in (0, \Gamma_\beta^*/(1 + \alpha))$, then we have

$$L_\delta^{\tilde{\epsilon}} < \frac{\log(C) + \log(1/\delta)}{\Gamma_\beta^* - (1 + \alpha)\epsilon} + 2 = \alpha \left[\frac{\log(C) + \log(1/\delta)}{\Gamma_\beta^*} \right] + 2 = \left[\frac{\alpha \log(C)}{\Gamma_\beta^*} + 2 \right] + \frac{\alpha \log(1/\delta)}{\Gamma_\beta^*}.$$

487 Therefore, we have

$$\limsup_{\delta \rightarrow 0} \frac{\mathbb{E}[\tau_\delta]}{\log(1/\delta)} \leq \limsup_{\delta \rightarrow 0} \frac{\mathbb{E}[M^\epsilon] + \mathbb{E}[L_\delta^{\tilde{\epsilon}}]}{\log(1/\delta)} \leq \frac{\alpha}{\Gamma_\beta^*}.$$

488 E Preliminaries

489 In this section, we introduce several preliminary results which is used in the theoretical analysis of
 490 TTEI.

491 **E.1 Properties of $f(x) = x\Phi(x) + \phi(x)$**

492 We provide several properties of the function $f(x) = x\Phi(x) + \phi(x)$ including its monotonicity, upper
493 bound and lower bound.

494 **Lemma 2.** $f(x)$ is positive and increasing on \mathbb{R} .

495 *Proof.* This is true since $f'(x) = \Phi(x) \geq 0$ and $\lim_{x \rightarrow -\infty} f(x) = 0$. □

496 **Lemma 3.** For $x > 0$,

$$f(-x) < \phi(-x).$$

497 *Proof.* For $x > 0$, $f(-x) = -x\Phi(-x) + \phi(-x) < \phi(-x)$. □

498 **Lemma 4.** For $x \geq 2$,

$$f(-x) > \frac{1}{x^3}\phi(-x).$$

499 *Proof.* Let $g(x) = \frac{1}{x}[f(-x) - \frac{1}{x^3}\phi(-x)] = -\Phi(-x) + \frac{1}{x}\phi(-x) - \frac{1}{x^4}\phi(-x)$. We have $g'(x) =$
500 $(-x^{-2} + x^{-3} + 4x^{-5})\phi(x) = x^{-5}(-x + 2)(x^2 + x + 2)\phi(x)$, which implies that $g(x)$ is decreasing
501 in $[2, \infty)$. We notice that $g(2) > 0$ and $\lim_{x \rightarrow \infty} g(x) = 0$, so for $x \geq 2$, $g(x) > 0$. Therefore, for
502 $x \geq 2$, $f(-x) > \frac{1}{x^3}\phi(-x)$. □

503 Lemmas 3 and 4 provides the upper and lower bounds for $f(\cdot)$, which is used to study the expected
504 improvement measures.

505 **E.2 Maximal Inequalities**

506 In the theoretical analysis of TTEI, we need a bound on the difference between the empirical mean
507 $\mu_{n,i}$ and the unknown true mean μ_i for each arm $i \in A$ at period n , and a bound on the difference
508 between $T_{n,i}$ and $\Psi_{n,i}$, two measurements of effort allocated to arm i up to period n . Two sample-path
509 dependent variables W_1 and W_2 are required to obtain the two bounds.

510 **Lemma 5.** Under any sampling rule beginning with an improper prior for each arm $i \in A$ with
511 $\mu_{1,i} = 0$ and $\sigma_{1,i} = \infty$, $\mathbb{E}[e^{\lambda W_1}] < \infty$ for all $\lambda > 0$ where

$$W_1 \triangleq \max_{n \in \mathbb{N}} \max_{i \in A} \sqrt{\frac{T_{n,i} + 1}{\log(e + T_{n,i})}} \left| \frac{\mu_{n,i} - \mu_i}{\sigma} \right|.$$

512 *Proof.* Under any sampling rule beginning with an improper prior for each arm $i \in A$ with $\sigma_{1,i} = \infty$
513 and $\mu_{1,i} = 0$ for each arm $i \in A$, if $T_{n,i} = \sum_{\ell=1}^{n-1} \mathbf{1}\{I_\ell = i\} = 0$, $\mu_{n,i} = \mu_{1,i} = 0$, and if $T_{n,i} > 0$,

$$\mu_{n,i} = \frac{1}{T_{n,i}} \sum_{\ell=1}^{n-1} \mathbf{1}\{I_\ell = i\} Y_{\ell, I_\ell}.$$

514 A mathematically equivalent way of simulating the system is to generate a collection of independent
515 variables $(X_{n,i})_{n \in \mathbb{N}, i \in A}$ where each $X_{n,i} \sim N(\mu_i, \sigma^2)$. At time n , the algorithm selects an arm I_n ,
516 and observes the real valued response X_{S_{n, I_n}, I_n} where $S_{n, I_n} \triangleq \sum_{\ell=1}^n \mathbf{1}\{I_\ell = i\}$. For all $i \in A$, we
517 let $\bar{X}_{0,i} = 0$, and for $n \in \mathbb{N}$, $\bar{X}_{n,i} = \frac{1}{n} \sum_{\ell=1}^n X_{\ell, i}$ denote the empirical mean of arm i up to the n th
518 time it is chosen. We will bound

$$\widetilde{W} \triangleq \max_{n \in \mathbb{N} \cup \{0\}} \max_{i \in A} \sqrt{\frac{n+1}{\log(e+n)}} \left| \frac{\bar{X}_{n,i} - \mu_i}{\sigma} \right|.$$

519 When every arm is played infinitely often, $W_1 = \widetilde{W}$. One always has $W_1 \leq \widetilde{W}$, so it is sufficient to
520 bound $\mathbb{E}[e^{\lambda \widetilde{W}}]$ for all $\lambda > 0$. Notice that $\widetilde{W} = \max\{\xi, |\mu_1|/\sigma, \dots, |\mu_k|/\sigma\} \leq \xi + \sigma^{-1} \sum_{i \in A} |\mu_i|$
521 where

$$\xi \triangleq \max_{n \in \mathbb{N}} \max_{i \in A} \sqrt{\frac{n+1}{\log(e+n)}} \left| \frac{\bar{X}_{n,i} - \mu_i}{\sigma} \right|.$$

522 Hence, it suffices to bound $\mathbb{E}[e^{\lambda\xi}]$ for all $\lambda > 0$.

523 For all $n \in \mathbb{N}$ and $i \in A$, we define $Z_{n,i} \triangleq \sqrt{n} \left(\frac{\bar{X}_{n,i} - \mu_i}{\sigma} \right)$, and then

$$\xi = \max_{n \in \mathbb{N}} \max_{i \in A} \sqrt{\frac{n+1}{n \log(e+n)}} |Z_{n,i}|.$$

524 Each $Z_{n,i} \sim N(0, 1)$, and thus by Lemma 1, $Z_{n,i}$ satisfies the tail bound $\mathbb{P}(|Z_{n,i}| \geq z) \leq e^{-z^2/2}$
 525 for $z > 0$. Therefore, for all $x \geq 2$

$$\begin{aligned} \mathbb{P}(\xi \geq 2x) &= \mathbb{P}\left(\exists n \in \mathbb{N}, i \in A : |Z_{n,i}| \geq 2\sqrt{\frac{n \log(e+n)}{n+1}} x\right) \\ &\leq \sum_{n,i} \mathbb{P}\left(|Z_{n,i}| \geq 2\sqrt{\frac{n \log(e+n)}{n+1}} x\right) \\ &\leq \sum_{n,i} \exp\left(-\frac{2n \log(e+n)}{n+1} x^2\right) \\ &= k \sum_n \exp\left(-\frac{2n \log(e+n)}{n+1} x^2\right) \\ &\stackrel{(*)}{\leq} k \sum_n \exp\left(-2 \log(e+n) - \frac{n}{n+1} x^2\right) \\ &= k \sum_n \left(\frac{1}{e+n}\right)^2 e^{-\frac{n}{n+1} x^2} \\ &\leq C e^{-x^2/2}. \end{aligned}$$

526 where step (*) uses the $ab \geq a + b$ when $a, b \geq 2$ and $C = k \sum_{n \in \mathbb{N}} (e+n)^{-2} < \infty$ is a constant.
 527 Then for all $\lambda > 0$,

$$\mathbb{E}[e^{\lambda\xi}] = \int_{x=1}^{\infty} \mathbb{P}(e^{\lambda\xi} \geq x) dx \stackrel{(*)}{=} \int_{u=0}^{\infty} \mathbb{P}(e^{\lambda\xi} \geq e^{2\lambda u}) 2\lambda e^{2\lambda u} du \leq 2 + C \int_{u=2}^{\infty} e^{-u^2/2} \cdot 2\lambda e^{2\lambda u} du < \infty$$

528 where in step (*), we have substituted $x = e^{2\lambda u}$. Hence, for all $\lambda > 0$, $\mathbb{E}[e^{\lambda W_1}] < \infty$. \square

529 This result provides a bound for the difference between the empirical mean of an arm and its true
 530 unknown mean. For $i \in A$ and $n \in \mathbb{N}$

$$|\mu_{n,i} - \mu_i| \leq \sigma W_1 \sqrt{\frac{\log(e + T_{n,i})}{T_{n,i} + 1}}.$$

531 Then we introduce the second sample-path dependent variable W_2 , and the following lemma on the
 532 difference between two measurements of effort under any top-two sampling rule, which at each time,
 533 measures one of the two designs that appear most promising given current evidence.

534 **Lemma 6.** *Under any top-two sampling rule with parameter $\beta \in (0, 1)$ beginning with an improper
 535 prior for each arm $i \in A$ with $\mu_{1,i} = 0$ and $\sigma_{1,i} = \infty$, $\mathbb{E}[e^{\lambda W_2}] < \infty$ for all $\lambda > 0$ where*

$$W_2 \triangleq \max_{n \in \mathbb{N}} \max_{i \in A} \frac{|T_{n,i} - \Psi_{n,i}|}{\sqrt{(1 + \Psi_{n,i}/\beta_{\min}) \log(e^2 + \Psi_{n,i}/\beta_{\min})}}.$$

536 *Proof.* Similar to the proof for Lemma 5, it suffices to show $\mathbb{P}(W_2 \geq x) \leq k e^{-x^2/2}$ for all $x \geq 2$.

537 Fix some $i \in A$. Define for each $n \in \mathbb{N}$

$$D_n \triangleq T_{n,i} - \Psi_{n,i} = \sum_{\ell=1}^{n-1} d_\ell$$

538 where

$$d_n \triangleq \mathbf{1}(I_n = i) - \psi_{n,i} = \mathbf{1}(I_n = i) - \mathbb{P}(I_n = i | \mathcal{F}_{n-1}).$$

539 Then $\mathbb{E}[d_n | \mathcal{F}_{n-1}] = 0$ and D_n is a zero mean martingale. Now, note $\psi_{n,i} \in \{0, \beta, 1 - \beta\}$ almost
540 surely, and set

$$X_n := \mathbf{1}(\psi_{n,i} > 0)$$

541 to be the indicator that i is among the top-two in period n . We can see that $d_n = X_n d_n$, and so

$$D_n = \sum_{\ell=1}^{n-1} X_\ell d_\ell.$$

542 Here $\{X_n\}$ is a binary valued previsible process (i.e. X_n is \mathcal{F}_{n-1} measurable), and d_n is a
543 zero-mean \mathcal{F}_n adapted process with increments bounded as $|d_n| \leq 1$ almost surely.

544 The quadratic variation of D_n is

$$\langle D \rangle_n = \sum_{\ell=1}^{n-1} \mathbb{E}[X_\ell d_\ell^2 | \mathcal{F}_{\ell-1}] = \sum_{\ell=1}^{n-1} X_\ell \beta (1 - \beta)$$

545 and so the magnitude of fluctuations of the martingale D_n scale with the number of times i is in the
546 top-two.

547 There are a number of martingale analogues to the central limit theorem, which suggest that $D_n =$
548 $O_P\left(\sqrt{\langle D \rangle_n}\right)$. To establish this formally, we apply the theorem of self-normalized martingale
549 processes [20], which bound processes like $D_n / \sqrt{\langle D \rangle_n}$. We will apply a result established in [1].

550 Because $|d_n| \leq 1$, applying Hoeffding's Lemma implies

$$E[e^{\lambda d_n} | \mathcal{F}_{n-1}] \leq e^{\lambda^2/2}, \quad \lambda \in \mathbb{R}$$

551 and so d_n is 1-sub-Gaussian conditioned on \mathcal{F}_{n-1} . Applying Corollary 8 of [1] implies that for any
552 $\delta > 0$, with probability least $1 - \delta$

$$|D_n| \leq \sqrt{2 \left(1 + \sum_{\ell=1}^{n-1} X_\ell\right) \log \left(\frac{\sqrt{1 + \sum_{\ell=1}^{n-1} X_\ell}}{\delta}\right)}, \quad \forall n \in \mathbb{N}$$

553 Analogously, for any $x \geq 2$ with probability at least $1 - e^{-x^2/2}$,

$$\begin{aligned} |D_n| &\leq \sqrt{2 \left(1 + \sum_{\ell=1}^{n-1} X_\ell\right) \log \left(\frac{\sqrt{1 + \sum_{\ell=1}^{n-1} X_\ell}}{e^{-x^2/2}}\right)} \\ &= \sqrt{\left(1 + \sum_{\ell=1}^{n-1} X_\ell\right) \left(\log \left(1 + \sum_{\ell=1}^{n-1} X_\ell\right) + x^2\right)} \\ &\leq \sqrt{\left(1 + \sum_{\ell=1}^{n-1} X_\ell\right) \left(\log \left(e^2 + \sum_{\ell=1}^{n-1} X_\ell\right) + x^2\right)} \\ &\leq \sqrt{\left(1 + \sum_{\ell=1}^{n-1} X_\ell\right) \log \left(e^2 + \sum_{\ell=1}^{n-1} X_\ell\right) x^2} \end{aligned}$$

554 for all $n \in \mathbb{N}$, where the last step uses that $ab \geq a + b$ for $a, b \geq 2$. Then, for all $x \geq 2$

$$\mathbb{P} \left(\max_{n \in \mathbb{N}} \frac{|D_n|}{\sqrt{\left(1 + \sum_{\ell=1}^{n-1} X_\ell\right) \log \left(e^2 + \sum_{\ell=1}^{n-1} X_\ell\right)}} \geq x \right) \leq e^{-x^2/2}$$

555 Since $\Psi_{n,i} \geq \beta_{\min} \sum_{\ell=1}^{n-1} X_\ell$, we have shown that for any i ,

$$\mathbb{P} \left(\max_{n \in \mathbb{N}} \frac{|T_{n,i} - \Psi_{n,i}|}{\sqrt{(1 + \Psi_{n,i}/\beta_{\min}) \log(e^2 + \Psi_{n,i}/\beta_{\min})}} \geq x \right) \leq e^{-x^2/2}$$

556 Taking a union bound over $i \in A$ implies $\mathbb{P}(W_2 \geq x) \leq ke^{-x^2/2}$ for any $x \geq 2$. \square

557 This result implies that for any period n and arm i ,

$$|T_{n,i} - \Psi_{n,i}| \leq W_2 \sqrt{(1 + \Psi_{n,i}/\beta_{\min}) \log(e^2 + \Psi_{n,i}/\beta_{\min})}.$$

558 The next result provides another bound, which is used in the theoretical analysis of TTEI.

559 **Lemma 7.** *Under TTEI with parameter $\beta \in (0, 1)$ beginning with an improper prior for each arm*
 560 *$i \in A$ with $\mu_{1,i} = 0$ and $\sigma_{1,i} = \infty$, for all $n \in \mathbb{N}$ and arm $i \in A$,*

$$|T_{n,i} - \Psi_{n,i}| < \left(2 + \frac{3\Psi_{n,i}^{3/4}}{\beta_{\min}} \right) W_2.$$

561 *Proof.* Fix some arm $i \in A$. If arm i is never chosen in either case 1 or case 2 of TTEI up to period
 562 n , then $\Psi_{n,i} = 0$, and thus

$$|T_{n,i} - \Psi_{n,i}| \leq W_2 \sqrt{(1 + \Psi_{n,i}/\beta_{\min}) \log(e^2 + \Psi_{n,i}/\beta_{\min})} < 2W_2$$

563 Once arm i has been chosen in either case 1 or case 2 of TTEI, $\Psi_{n,i} \geq \beta_{\min}$. Then we have
 564 $1 + \Psi_{n,i}/\beta_{\min} < 3\Psi_{n,i}/\beta_{\min}$ and $\log(e^2 + \Psi_{n,i}/\beta_{\min}) < 3(\Psi_{n,i}/\beta_{\min})^{1/2}$, which leads to

$$|T_{n,i} - \Psi_{n,i}| < 3W_2(\Psi_{n,i}/\beta_{\min})^{3/4} < \frac{3\Psi_{n,i}^{3/4}}{\beta_{\min}} W_2.$$

565 Hence,

$$|T_{n,i} - \Psi_{n,i}| < \max \left\{ 2, \frac{3\Psi_{n,i}^{3/4}}{\beta_{\min}} \right\} W_2 < \left(2 + \frac{3\Psi_{n,i}^{3/4}}{\beta_{\min}} \right) W_2.$$

566 \square

567 E.3 Technical Lemmas

568 The following technical lemma is used to quantify the time after which TTEI satisfies a certain
 569 property. We want to write such a time as a polynomial of sample-path dependent variables.

570 **Lemma 8.** *Fix constants $c_0 > c_1 > 0$ and $c, c_2 > 0$. Then for any $a_1, a_2 > 0$, there exists a*
 571 *$X = \text{poly}(a_1, a_2)$ such that for all $x \geq X$,*

$$\exp(cx^{c_0} - a_1x^{c_1}) > a_2x^{c_2}.$$

572 *Proof.* There exists $X_1 = \text{poly}(a_1)$ such that for all $x \geq X_1$, $cx^{c_0-c_1} - a_1 > 1$. In addition, there
 573 exists $X_2 = \text{poly}(a_2)$ such that for all $x \geq X_2$, $\exp(x^{c_1}) > a_2x^{c_2}$. Hence, for all $x \geq X \triangleq$
 574 $\max\{X_1, X_2\}$,

$$\exp(cx^{c_0} - a_1x^{c_1}) = \exp(x^{c_1} (cx^{c_0-c_1} - a_1)) \geq \exp(x^{c_1}) > a_2x^{c_2}.$$

575 \square

576 F Results specific to TTEI

577 In this section, we present theoretical results specific to the proposed TTEI policy. The main challenge
 578 is ensuring $\mathbb{E}[T_\beta^\epsilon]$ is finite where T_β^ϵ is the time after which for each arm, its empirical mean and
 579 empirical proportion are ϵ -accurate. To do this, we present several results for any sample path (up
 580 to a set of measure zero), and show that T_β^ϵ depends at most polynomially on W_1 and W_2 . By
 581 Lemmas 5 and 6, the expected value of polynomials of W_1 and W_2 is finite. This ensures that $\mathbb{E}[T_\beta^\epsilon]$
 582 is finite, which immediately establishes that TTEI achieves the sufficient conditions for both notions
 583 of optimality.

584 **F1 Sufficient Exploration**

585 We first show that every arm is sampled frequently under TTEI.

586 **Proposition 2.** *Under TTEI with parameter $\beta \in (0, 1)$, there exists $N_1 = \text{poly}(W_1, W_2)$ such that*
 587 *for all $n \geq N_1$,*

$$T_{n,i} \geq \sqrt{n/k}, \quad \forall i \in A.$$

588 To prove this proposition, we first need to define two under-sampled sets for all $L > 0$ and $n \in \mathbb{N}$:

$$U_n^L \triangleq \{i \in A : T_{n,i} < L^{1/2}\}$$

589 and

$$V_n^L \triangleq \{i \in A : T_{n,i} < L^{3/4}\}.$$
¹

590 Let $\overline{U_n^L} \triangleq A \setminus U_n^L$ and $\overline{V_n^L} \triangleq A \setminus V_n^L$. Then Proposition 2 can be proved using the following two
 591 lemmas. Note that in this paper, $X = \text{poly}(W_1, W_2)$ means that $X = \mathcal{O}(W_1^{c_1} W_2^{c_2})$ for positive
 592 constants c_1 and c_2 where $(\sigma, k, \mu_1, \dots, \mu_k, \beta)$ are treated as constants throughout the proof.

593 **Lemma 9.** *Under TTEI with parameter $\beta \in (0, 1)$, there exists $L_1 = \text{poly}(W_1)$ such that for all*
 594 *$L \geq L_1$ and $n \leq kL$,² if U_n^L is nonempty, then $I_n^{(1)} \in V_n^L$ or $I_n^{(2)} \in V_n^L$.*

595 *Proof.* First of all, we will show that if $I_n^{(1)} \in \overline{V_n^L}$, then $I_n^* \in \overline{V_n^L}$ where $I_n^* = \arg \max_{i \in A} \mu_{n,i}$.
 596 We prove this by contradiction. Suppose $I_n^* \in V_n^L$. By definition, $T_{n, I_n^{(1)}} > T_{n, I_n^*}$, which implies
 597 $\sigma_{n, I_n^{(1)}} < \sigma_{n, I_n^*}$. By Lemma 2, we have

$$v_{n, I_n^{(1)}}^{(1)} = \sigma_{n, I_n^{(1)}} f\left(\frac{\mu_{n, I_n^{(1)}} - \mu_{n, I_n^*}}{\sigma_{n, I_n^{(1)}}}\right) < \sigma_{n, I_n^*} f(0) = v_{n, I_n^*}^{(1)},$$

598 which contradicts the definition of $I_n^{(1)}$. Hence, if $I_n^{(1)} \in \overline{V_n^L}$, then $I_n^* \in \overline{V_n^L}$.

599 Secondly we will show that when L is sufficiently large, if $I_n^* \in \overline{V_n^L}$, then for all $i \in \overline{V_n^L} \setminus \{I_n^*\}$,
 600 $\mu_{n,i} - \mu_{n, I_n^*} \leq -0.5\Delta_{\min}$ where $\Delta_{\min} = \min_{i \neq j} |\mu_i - \mu_j| > 0$. By Lemma 5, for all $i \in \overline{V_n^L}$,

$$|\mu_{n,i} - \mu_i| \leq \sigma W_1 \sqrt{\frac{\log(e + T_{n,i})}{T_{n,i} + 1}} \leq \sigma W_1 \sqrt{\frac{\log(e + L^{3/4})}{L^{3/4} + 1}}$$

601 where the last inequality is valid because $g(x) = \log(e + x)/(x + 1)$ is positive and decreasing
 602 on $(0, \infty)$ and $T_{n,i} \geq L^{3/4}$. Note that for $L \geq 1$, $\log(e + L^{3/4}) \leq 2L^{1/4}$. Then there exists
 603 $M_1 = \text{poly}(W_1)$ such that for all $L \geq M_1$,

$$\sqrt{\frac{\log(e + L^{3/4})}{L^{3/4} + 1}} \leq \sqrt{\frac{2L^{1/4}}{L^{3/4} + 1}} \leq \frac{\Delta_{\min}}{4\sigma W_1}.$$

604 Suppose there exists $\tilde{i} \in \overline{V_n^L} \setminus \{I_n^*\}$ such that $\mu_{\tilde{i}} > \mu_{I_n^*}$. Then for $L \geq M_1$, we have

$$\begin{aligned} \mu_{n, \tilde{i}} - \mu_{n, I_n^*} &\geq \mu_{\tilde{i}} - \sigma W_1 \sqrt{\frac{\log(e + L^{3/4})}{L^{3/4} + 1}} - \mu_{I_n^*} - \sigma W_1 \sqrt{\frac{\log(e + L^{3/4})}{L^{3/4} + 1}} \\ &= (\mu_{\tilde{i}} - \mu_{I_n^*}) - 2\sigma W_1 \sqrt{\frac{\log(e + L^{3/4})}{L^{3/4} + 1}} \\ &\geq \Delta_{\min} - 2\sigma W_1 (\Delta_{\min}/4\sigma W_1) = 0.5\Delta_{\min}, \end{aligned}$$

¹We fix the exponent here to be $3/4$. Indeed, it can be changed to $1/2 + \epsilon$ for any $\epsilon > 0$. We just need a gap between the exponent here and $1/2$ in U_n^L .

² L could be any value, but n must be integer value.

605 which contradicts the definition of I_n^* . Hence, for $L \geq M_1$, if $I_n^* \in \overline{V_n^L}$, then $\mu_{I_n^*} > \mu_i$ for all
 606 $i \in \overline{V_n^L} \setminus \{I_n^*\}$ (note that we assume that all arm-means are unique), and thus

$$\mu_{n,i} - \mu_{n,I_n^*} \leq (\mu_i - \mu_{I_n^*}) + 2\sigma W_1 \sqrt{\frac{\log(e + L^{3/4})}{L^{3/4} + 1}} \leq -\Delta_{\min} + 0.5\Delta_{\min} = -0.5\Delta_{\min}.$$

607 Thirdly we will show when L is sufficiently large and $n \leq kL$, if $I_n^{(1)} \in \overline{V_n^L}$ (which implies
 608 $I_n^* \in \overline{V_n^L}$), then $v_{n,I_n^*}^{(1)} > v_{n,i}^{(1)}$ for all $i \in \overline{V_n^L} \setminus \{I_n^*\}$, which implies $I_n^{(1)} = I_n^*$. For all $i \in \overline{V_n^L} \setminus \{I_n^*\}$,
 609 $\sigma_{n,i}^2 = \sigma^2/T_{n,i} \leq \sigma^2/L^{3/4}$, and when $L \geq M_1$, $\mu_{n,i} - \mu_{n,I_n^*} \leq -0.5\Delta_{\min}$, which lead to

$$v_{n,i}^{(1)} = \sigma_{n,i} f\left(\frac{\mu_{n,i} - \mu_{n,I_n^*}}{\sigma_{n,i}}\right) \leq \frac{\sigma}{L^{3/8}} f\left(\frac{-\Delta_{\min} L^{3/8}}{2\sigma}\right) < \frac{\sigma}{L^{3/8}} \phi\left(\frac{-\Delta_{\min} L^{3/8}}{2\sigma}\right) \quad (5)$$

610 where the last inequality uses Lemma 3. On the other hand,

$$v_{n,I_n^*}^{(1)} = \sigma_{n,I_n^*} f(0) \geq \frac{\sigma}{(kL)^{1/2}} \phi(0). \quad (6)$$

611 There exists M_2 such that for all $L \geq M_2$, the right hand side of (6) is larger than the right hand
 612 of (5). Hence, for $L \geq \max\{M_1, M_2\}$ and $n \leq kL$, if $I_n^{(1)} \in \overline{V_n^L}$ (which implies $I_n^* \in \overline{V_n^L}$), then
 613 $v_{n,I_n^*}^{(1)} > v_{n,i}^{(1)}$ for all $i \in \overline{V_n^L} \setminus \{I_n^*\}$, which implies $I_n^{(1)} = I_n^*$.

614 Finally we will show that when L is sufficiently large and $n \leq kL$, if U_n^L is nonempty (which implies
 615 V_n^L is nonempty by definition) and $I_n^{(1)} \in \overline{V_n^L}$ (which implies $I_n^* \in \overline{V_n^L}$), then $I_n^{(2)} \in V_n^L$. We have
 616 proved that for $L \geq \{M_1, M_2\}$, $I_n^{(1)} = I_n^*$. Then for all $i \in \overline{V_n^L} \setminus \{I_n^*\}$,

$$\mu_{n,i} - \mu_{n,I_n^{(1)}} = \mu_{n,i} - \mu_{n,I_n^*} \leq -0.5\Delta_{\min},$$

617 and by definition,

$$\sigma_{n,i}^2 + \sigma_{n,I_n^{(1)}}^2 = \sigma_{n,i}^2 + \sigma_{n,I_n^*}^2 = \frac{\sigma^2}{T_{n,i}} + \frac{\sigma^2}{T_{n,I_n^*}} \leq \frac{\sigma^2}{L^{3/4}} + \frac{\sigma^2}{L^{3/4}} < \frac{4\sigma^2}{L^{3/4}},$$

618 which leads to

$$v_{n,i}^{(2)} < \frac{2\sigma}{L^{3/8}} f\left(\frac{-\Delta_{\min} L^{3/8}}{4\sigma}\right) < \frac{2\sigma}{L^{3/8}} \phi\left(\frac{-\Delta_{\min} L^{3/8}}{4\sigma}\right). \quad (7)$$

619 where the last inequality uses Lemma 3. On the other hand, for all $j \in U_n^L$,

$$\begin{aligned} \mu_{n,j} - \mu_{n,I_n^{(1)}} &= \mu_{n,j} - \mu_{n,I_n^*} \\ &\geq \mu_j - \sigma W_1 \sqrt{\frac{\log(e + T_{n,j})}{T_{n,j} + 1}} - \mu_{I_n^*} - \sigma W_1 \sqrt{\frac{\log(e + T_{n,I_n^*})}{T_{n,I_n^*} + 1}} \\ &\geq (\mu_j - \mu_{I_n^*}) - 2\sigma W_1 \sqrt{\frac{\log(e)}{1}} = (\mu_j - \mu_{I_n^*}) - 2\sigma W_1 \end{aligned}$$

620 where the last inequality is valid because $g(x) = \log(e + x)/(x + 1)$ is positive and decreasing
 621 on $(0, \infty)$ and $T_{n,j}, T_{n,I_n^*} \geq 0$. If $\mu_{I_n^*} > \mu_j$, $\mu_{n,j} - \mu_{n,I_n^{(1)}} \geq -\Delta_{\max} - 2\sigma W_1$ where $\Delta_{\max} =$
 622 $\max_{i,j \in A} (\mu_i - \mu_j)$; otherwise, $\mu_{n,j} - \mu_{n,I_n^{(1)}} \geq \Delta_{\min} - 2\sigma W_1 > -\Delta_{\max} - 2\sigma W_1$. Hence, we
 623 have $\mu_{n,j} - \mu_{n,I_n^{(1)}} \geq -\Delta_{\max} - 2\sigma W_1$, and by definition,

$$\sigma_{n,j}^2 + \sigma_{n,I_n^{(1)}}^2 = \sigma_{n,j}^2 + \sigma_{n,I_n^*}^2 = \frac{\sigma^2}{T_{n,j}} + \frac{\sigma^2}{T_{n,I_n^*}} > \frac{\sigma^2}{L^{1/2}} + \frac{\sigma^2}{kL} > \frac{\sigma^2}{L^{1/2}},$$

624 which leads to

$$v_{n,j}^{(2)} > \frac{\sigma}{L^{1/4}} f\left(\frac{-(\Delta_{\max} + 2\sigma W_1)L^{1/4}}{\sigma}\right).$$

625 Let $M_3 \triangleq (2\sigma/\Delta_{\max})^4$. Since $W_1 \geq 0$ by definition, for all $L \geq M_3$, $(\Delta_{\max} + 2\sigma W_1)L^{1/4}/\sigma \geq 2$,
 626 and then by Lemma 4, we have

$$v_{n,j}^{(2)} > \frac{\sigma}{L^{1/4}} f\left(\frac{-(\Delta_{\max} + 2\sigma W_1)L^{1/4}}{\sigma}\right) > \frac{\sigma^4}{L(\Delta_{\max} + 2\sigma W_1)^3} \phi\left(\frac{-(\Delta_{\max} + 2\sigma W_1)L^{1/4}}{\sigma}\right). \quad (8)$$

627 By Lemma 8, there exists M_4 such that for all $L \geq M_4$, the right hand side of (8) is larger than
 628 the right hand side of (7). Therefore, for $L \geq L_1 \triangleq \max\{M_1, M_2, M_3, M_4\}$ and $n \leq kL$, if U_n^L is
 629 nonempty (which implies V_n^L is nonempty by definition) and $I_n^{(1)} \in \overline{V_n^L}$ (which implies $I_n^* \in \overline{V_n^L}$),
 630 then $v_{n,j}^{(2)} > v_{n,i}^{(2)}$ for all $j \in U_n^L$ and $i \in \overline{V_n^L}$ (here we use $v_{n,I_n^*}^{(2)} = v_{n,I_n^{(1)}}^{(2)} = 0$), which implies
 631 $I_n^{(2)} \notin \overline{V_n^L}$, and thus $I_n^{(2)} \in V_n^L$.

632 □

633 Note that the floor function $\lfloor x \rfloor$ is the greatest integer less than or equal to x . Then based on Lemma
 634 9, we have the following result.

635 **Lemma 10.** *Under TTEI with parameter $\beta \in (0, 1)$, there exists $L_2 = \text{poly}(W_1, W_2)$ such that for
 636 all $L \geq L_2$, $U_{\lfloor kL \rfloor}^L$ is empty.*

637 *Proof.* There exists $M_1 = \text{poly}(W_2)$ such that for all $L \geq M_1$, we have $\lfloor L \rfloor - 1 \geq kL^{3/4}$ and

$$\beta_{\min} \lfloor L \rfloor - 4kW_2 - \frac{6k \lfloor kL \rfloor^{3/4}}{\beta_{\min}} W_2 \geq kL^{3/4}$$

638 where $\beta_{\min} = \min\{\beta, 1 - \beta\} > 0$. Let $L_2 \triangleq \max\{L_1, M_1\}$ where $L_1 = \text{poly}(W_1)$ has been
 639 introduced in Lemma 9. Now We want to prove this statement by contradiction.

640 Suppose there exists some $L \geq L_2$ such that $U_{\lfloor kL \rfloor}^L$ is nonempty. Then all
 641 $U_1^L, U_2^L, \dots, U_{\lfloor kL \rfloor - 1}^L, U_{\lfloor kL \rfloor}^L$ are nonempty, and thus by definition, all $V_1^L, V_2^L, \dots, V_{\lfloor kL \rfloor - 1}^L, V_{\lfloor kL \rfloor}^L$
 642 are empty. Since $L \geq L_2$, we have $\lfloor L \rfloor - 1 \geq kL^{3/4}$, so at least one arm is measured at least $L^{3/4}$
 643 times before period $\lfloor L \rfloor$, and thus $|V_{\lfloor L \rfloor}^L| \leq k - 1$.

644 Now we want to prove $|V_{\lfloor 2L \rfloor}^L| \leq k - 2$. For all $\ell = \lfloor L \rfloor, \lfloor L \rfloor + 1, \dots, \lfloor 2L \rfloor - 1$, U_ℓ^L is nonempty,
 645 then by Lemma 9, we have $I_n^{(1)} \in V_\ell^L$ or $I_n^{(2)} \in V_\ell^L$, and thus $\sum_{i \in V_\ell^L} \psi_{\ell,i} = \sum_{i \in V_\ell^L} \mathbb{P}(I_\ell =$
 646 $i | \mathcal{F}_{\ell-1}) \geq \beta_{\min}$, which implies $\sum_{i \in V_{\lfloor L \rfloor}^L} \psi_{\ell,i} \geq \beta_{\min}$ due to $V_\ell^L \subseteq V_{\lfloor L \rfloor}^L$. Hence, we have

$$\sum_{i \in V_{\lfloor L \rfloor}^L} (\Psi_{\lfloor 2L \rfloor, i} - \Psi_{\lfloor L \rfloor, i}) = \sum_{\ell = \lfloor L \rfloor}^{\lfloor 2L \rfloor - 1} \sum_{i \in V_{\lfloor L \rfloor}^L} \psi_{\ell, i} \geq \beta_{\min} \lfloor L \rfloor$$

647 where the inequality uses the fact that $\lfloor a + b \rfloor \geq \lfloor a \rfloor + \lfloor b \rfloor$ for $a, b \geq 0$. Then by Lemma 7, we have

$$\begin{aligned} & \sum_{i \in V_{\lfloor L \rfloor}^L} (T_{\lfloor 2L \rfloor, i} - T_{\lfloor L \rfloor, i}) \\ & \geq \sum_{i \in V_{\lfloor L \rfloor}^L} (\Psi_{\lfloor 2L \rfloor, i} - \Psi_{\lfloor L \rfloor, i}) - \sum_{i \in V_{\lfloor L \rfloor}^L} \left[\left(2 + \frac{3\Psi_{\lfloor 2L \rfloor, i}^{3/4}}{\beta_{\min}} \right) W_2 + \left(2 + \frac{3\Psi_{\lfloor L \rfloor, i}^{3/4}}{\beta_{\min}} \right) W_2 \right] \\ & \geq \beta_{\min} \lfloor L \rfloor - 2 \sum_{i \in V_{\lfloor L \rfloor}^L} \left(2 + \frac{3\Psi_{\lfloor kL \rfloor, i}^{3/4}}{\beta_{\min}} \right) W_2 \\ & > \beta_{\min} \lfloor L \rfloor - 2k \left(2 + \frac{3\Psi_{\lfloor kL \rfloor, i}^{3/4}}{\beta_{\min}} \right) W_2 \\ & > \beta_{\min} \lfloor L \rfloor - 4kW_2 - \frac{6k \lfloor kL \rfloor^{3/4}}{\beta_{\min}} W_2 \geq kL^{3/4} \end{aligned}$$

648 where the second last inequality uses that for all $i \in A$ and $n \in \mathbb{N}$, $\Psi_{n,i} \leq \beta_{\max}(n-1) < n$, and
 649 the last inequality is valid because of the construction of L_2 and $L \geq L_2$. Hence, at least one arm in
 650 $V_{\lfloor L \rfloor}^L$ is measured at least $L^{3/4}$ times in periods $[\lfloor L \rfloor, \lfloor 2L \rfloor]$, and thus $|V_{\lfloor 2L \rfloor}^L| \leq k-2$.

651 Similarly, we can prove that for $r = 3, \dots, k$, at least one arm in $V_{\lfloor (r-1)L \rfloor}^L$ is measured at least $L^{3/4}$
 652 times in periods $[\lfloor (r-1)L \rfloor, \lfloor rL \rfloor]$, so $|V_{\lfloor rL \rfloor}^L| \leq k-r$. Hence, $|V_{\lfloor kL \rfloor}^L| = 0$, i.e., $V_{\lfloor kL \rfloor}^L$ is empty,
 653 which implies that $U_{\lfloor kL \rfloor}^L$ is empty. \square

654 Now we can prove Proposition 2.

655 **Proof of Proposition 2.** Let $N_1 = kL_2$ where $L_2 = \text{poly}(W_1, W_2)$ introduced in Lemma 10. For
 656 all $n \geq N_1$, we let $L = n/k$, then by Lemma 10, we have $U_{\lfloor kL \rfloor}^L = U_n^{n/k}$ is empty, which by
 657 definition results in that for all $i \in A$, $T_{n,i} \geq \sqrt{n/k}$.

658 F.2 Concentration of Empirical Means

659 When n is large, using the bound on the difference between the empirical mean $\mu_{n,i}$ and the unknown
 660 true mean μ_i in terms of $T_{n,i}$ for each arm $i \in A$, we can formally show the concentration of $\mu_{n,i}$ to
 661 μ_i under TTEI.

662 **Proposition 3.** Let $\epsilon > 0$. Under TTEI with parameter $\beta \in (0, 1)$, there exists $N_2^\epsilon =$
 663 $\text{poly}(W_1, W_2, 1/\epsilon)$ such that for all $n \geq N_2^\epsilon$,

$$|\mu_{n,i} - \mu_i| \leq \epsilon, \quad \forall i \in A.$$

664 *Proof.* By Lemma 5, for all $i \in A$ and $n \in \mathbb{N}$,

$$|\mu_{n,i} - \mu_i| \leq \sigma W_1 \sqrt{\frac{\log(e + T_{n,i})}{T_{n,i} + 1}}.$$

665 By Proposition 2, for all $n \geq N_1$, for all $i \in A$, $T_{n,i} \geq \sqrt{n/k}$, and thus

$$|\mu_{n,i} - \mu_i| \leq \sigma W_1 \sqrt{\frac{\log(e + T_{n,i})}{T_{n,i} + 1}} \leq \sigma W_1 \sqrt{\frac{\log(e + (n/k)^{1/2})}{(n/k)^{1/2} + 1}}$$

666 where the last inequality uses $g(x) = \log(e+x)/(x+1)$ is positive and decreasing on $(0, \infty)$. Note
 667 that for $n \geq k$, $\log(e + (n/k)^{1/2}) \leq 2(n/k)^{1/4}$. Then there exists $M_1^\epsilon = \text{poly}(W_1, 1/\epsilon)$ such that
 668 for all $n \geq M_1^\epsilon$,

$$\sqrt{\frac{\log(e + (n/k)^{1/2})}{(n/k)^{1/2} + 1}} \leq \sqrt{\frac{2(n/k)^{1/4}}{(n/k)^{1/2} + 1}} \leq \frac{\epsilon}{\sigma W_1}.$$

669 Then for all $i \in A$ and $n \geq N_2^\epsilon \triangleq \max\{N_1, k, M_1^\epsilon\}$ where $N_1 = \text{poly}(W_1, W_2)$ introduced in
 670 Proposition 2, we have $|\mu_{n,i} - \mu_i| \leq \sigma W_1 [\epsilon/(\sigma W_1)] = \epsilon$. \square

671 Recall that we assume the unknown arm-means are unique and $\mu_1 > \mu_2 \dots > \mu_k$. If we set ϵ to a
 672 very small value in Lemma 3, when n is large, the empirical means are order as the true means, i.e.,
 673 $\mu_{n,1} > \mu_{n,2} \dots > \mu_{n,k}$, which implies the arm with the largest empirical mean is arm 1. In addition,
 674 we show that when n is large, the arm selected in case 1 of TTEI is also arm 1.

675 **Lemma 11.** Under TTEI with parameter $\beta \in (0, 1)$, there exists $N_3 = \text{poly}(W_1, W_2)$ such that for
 676 all $n \geq N_3$, $I_n^{(1)} = I_n^* = 1$.

677 *Proof.* Let $M_1 \triangleq N_2^{\Delta_{\min}/4}$. By Proposition 3, for all $n \geq M_1$,

$$|\mu_{n,i} - \mu_i| \leq \Delta_{\min}/4, \quad \forall i \in A$$

678 where $\Delta_{\min} = \min_{i \neq j} |\mu_i - \mu_j| > 0$, which implies $\mu_{n,1} > \mu_{n,2} > \dots > \mu_{n,k}$, and thus $I_n^* = 1$.

679 Now for $n \geq M_1$ and $i \neq I_n^*$, we have

$$\begin{aligned} \mu_{n,I_n^*} - \mu_{n,i} &= \mu_{n,1} - \mu_{n,i} \\ &\geq \mu_1 - \Delta_{\min}/4 - \mu_i - \Delta_{\min}/4 \\ &= (\mu_1 - \mu_i) - \Delta_{\min}/2 \\ &\geq \Delta_{\min} - \Delta_{\min}/2 = \Delta_{\min}/2. \end{aligned}$$

680 By Proposition 2, for $n \geq N_1$, $T_{n,i} \geq \sqrt{n/k}$ for all $i \in A$. Hence, for $n \geq \max\{N_1, M_1\}$ and
681 $i \neq I_n^*$, we have

$$v_{n,i}^{(1)} = \sigma_{n,i} f\left(\frac{\mu_{n,i} - \mu_{n,I_n^*}}{\sigma_{n,i}}\right) \leq \frac{\sigma k^{1/4}}{n^{1/4}} f\left(\frac{-\Delta_{\min} n^{1/4}}{2\sigma k^{1/4}}\right) < \frac{\sigma k^{1/4}}{n^{1/4}} \phi\left(\frac{-\Delta_{\min} n^{1/4}}{2\sigma k^{1/4}}\right) \quad (9)$$

682 where the two inequalities use Lemmas 2 and 3, respectively. On the other hand,

$$v_{n,I_n^*}^{(1)} = \sigma_{n,I_n^*} f(0) = \sigma_{n,I_n^*} \phi(0) > \frac{\sigma}{n^{1/2}} \phi(0) \quad (10)$$

683 where the inequality uses $T_{n,I_n^*} \leq n-1 < n$. There exists M_2 such that for all $n \geq M_2$, the right hand
684 side of (10) is larger than the right hand side of (9). Hence, for all $n \geq N_3 \triangleq \max\{N_1, M_2, M_2\}$,
685 $v_{n,I_n^*}^{(1)} > v_{n,i}^{(1)}$ for all $i \neq I_n^*$, which implies $I_n^{(1)} = I_n^* = 1$. \square

686 F.3 Tracking the Asymptotic Proportion of the Best Arm

687 In this subsection, we show that when the number of arm draws goes large, the empirical proportion
688 for the best arm concentrates to the tuning parameter β used in TTEI.

689 **Lemma 12.** *Let $\epsilon > 0$. Under TTEI with parameter $\beta \in (0, 1)$, there exists $N_4^\epsilon = \text{poly}(W_1, W_2, 1/\epsilon)$
690 such that for all $n \geq N_4^\epsilon$,*

$$\left| \frac{\Psi_{n,1}}{n} - \beta \right| \leq \epsilon.$$

691 *Proof.* By Lemma 11, for all $n \geq N_3$, we have $I_n^{(1)} = 1$. Then we have

$$\begin{aligned} \frac{\Psi_{n,1}}{n} &= \frac{1}{n} \left(\sum_{\ell=1}^{N_3-1} \psi_{\ell,1} + \sum_{\ell=N_3}^{n-1} \psi_{\ell,1} \right) \\ &\leq \frac{1}{n} [\beta_{\max}(N_3 - 1) + \beta(n - N_3)] \\ &< \beta + \frac{(\beta_{\max} - \beta)N_3}{n} \end{aligned}$$

692 where $\beta_{\max} = \max\{\beta, 1 - \beta\}$, and

$$\begin{aligned} \frac{\Psi_{n,1}}{n} &= \frac{1}{n} \left(\sum_{\ell=1}^{N_3-1} \psi_{\ell,1} + \sum_{\ell=N_3}^{n-1} \psi_{\ell,1} \right) \\ &\geq \frac{1}{n} \beta(n - N_3) \\ &= \beta - \frac{\beta N_3}{n}. \end{aligned}$$

693 For all $n \geq \beta_{\max} N_3 / \epsilon$, we have $(\beta_{\max} - \beta)N_3/n < \epsilon$ and $-\beta N_3/n \geq -\epsilon$. Therefore, for all
694 $n \geq N_4^\epsilon \triangleq \max\{N_3, \beta_{\max} N_3 / \epsilon\}$, we have $|\Psi_{n,1}/n - \beta| \leq \epsilon$. \square

695 Based on Lemma 12, we can prove the next result showing the concentration of $T_{n,1}/n$ to β .

696 **Lemma 13.** *Let $\epsilon > 0$. Under TTEI with parameter $\beta \in (0, 1)$, there exists $N_5^\epsilon = \text{poly}(W_1, W_2, 1/\epsilon)$
697 such that for all $n \geq N_5^\epsilon$,*

$$\left| \frac{T_{n,1}}{n} - \beta \right| \leq \epsilon.$$

698 *Proof.* It suffices to prove this statement for $\epsilon \in (0, \beta)$. By Lemma 12, for all $n \geq N_4^{\epsilon/2}$,
 699 $|\Psi_{n,1}/n - \beta| \leq \epsilon/2$, which implies $\Psi_{n,1} \geq (\beta - \epsilon/2)n$. Lemma 7 implies that for all
 700 $n \geq M_1^\epsilon \triangleq \max\{N_4^{\epsilon/2}, 2/\beta\}$,

$$\left| \frac{T_{n,1}}{\Psi_{n,1}} - 1 \right| \leq \left(\frac{2}{\Psi_{n,1}^{1/4}} + \frac{3}{\beta_{\min} \Psi_{n,1}^{1/4}} \right) W_2 \leq \frac{(2 + 3/\beta_{\min})W_2}{(\beta - \epsilon/2)^{1/4} n^{1/4}} < \frac{(2 + 3/\beta_{\min})W_2}{(\beta/2)^{1/4} n^{1/4}} \quad (11)$$

701 where the second inequality is valid since $\Psi_{n,1} \geq (\beta - \epsilon/2)n > (\beta/2)n \geq 1$. There exists
 702 $M_2^\epsilon = \text{poly}(W_2, 1/\epsilon)$ such that for all $n \geq M_2^\epsilon$, the right hand side of (11) is less than $\epsilon/(2\beta + \epsilon)$.
 703 Hence, for all $n \geq N_5^\epsilon \triangleq \max\{M_1^\epsilon, M_2^\epsilon\}$, $|T_{n,1}/\Psi_{n,1} - 1| < \epsilon/(2\beta + \epsilon)$ and $|\Psi_{n,1}/n - \beta| \leq \epsilon/2$,
 704 and thus we have

$$\frac{T_{n,1}}{n} < \left(1 + \frac{\epsilon}{2\beta + \epsilon} \right) \frac{\Psi_{n,1}}{n} \leq \left(1 + \frac{\epsilon}{2\beta + \epsilon} \right) (\beta + \epsilon/2) = \beta + \epsilon$$

705 and

$$\frac{T_{n,1}}{n} > \left(1 - \frac{\epsilon}{2\beta + \epsilon} \right) \frac{\Psi_{n,1}}{n} \geq \left(1 - \frac{\epsilon}{2\beta + \epsilon} \right) (\beta - \epsilon/2) > \beta - \epsilon,$$

706 which leads to $|T_{n,1}/n - \beta| < \epsilon$. \square

707 F.4 Tracking the Asymptotic Proportions of All Arms

708 Besides the best arm, we can further show that for each arm, its empirical proportion concentrates to
 709 its optimal proportion when the number of arm draws goes large.

710 **Proposition 4.** *Let $\epsilon > 0$. Under TTEI with parameter $\beta \in (0, 1)$, there exists $N_7^\epsilon =$
 711 $\text{poly}(W_1, W_2, 1/\epsilon, \epsilon)$ such that for all $n \geq N_7^\epsilon$,*

$$\left| \frac{T_{n,i}}{n} - w_i^\beta \right| \leq \epsilon, \quad \forall i \in A.$$

712 To prove this proposition, we need some further notations. For any $n \in \mathbb{N}$, we define the under-
 713 sampled set

$$P_n = \left\{ i \neq 1 : \frac{T_{n,i}}{n} - w_i^\beta < 0 \right\},$$

714 where the unique vector $(w_2^\beta, \dots, w_k^\beta)$ satisfies $\sum_{i=2}^k w_i^\beta = 1 - \beta$ and

$$\frac{(\mu_2 - \mu_1)^2}{1/w_2^\beta + 1/\beta} = \dots = \frac{(\mu_k - \mu_1)^2}{1/w_k^\beta + 1/\beta}.$$

715 Then given $\epsilon > 0$, we define the over-sampled set

$$O_n^\epsilon = \left\{ i \neq 1 : \frac{T_{n,i}}{n} - w_i^\beta > \epsilon \right\}.$$

716 The next result shows that when n is large, the over-sampled set is empty. Based on this result, we
 717 can prove that when n is large, the under-sampled set is also empty, which immediately establishes
 718 Proposition 4.

719 **Lemma 14.** *Let $\epsilon > 0$. Under TTEI with parameter $\beta \in (0, 1)$, there exists $N_6^\epsilon =$
 720 $\text{poly}(W_1, W_2, 1/\epsilon, \epsilon)$ such that for all $n \geq N_6^\epsilon$, O_n^ϵ is empty.*

721 *Proof.* If $O_n^{\epsilon/2}$ is empty, then O_n^ϵ is empty. Now let us consider the case that $O_n^{\epsilon/2}$ is nonempty, and
 722 it suffices to prove the statement for $\epsilon \in (0, \min\{\Delta_{\min}/2, 1\})$.

723 Fix $\epsilon \in (0, \min\{\Delta_{\min}/2, 1\})$. For $\epsilon' \in (0, \epsilon/2)$, by Proposition 3 and Lemma 13, we can find large
 724 enough $M_1^{\epsilon'} = \text{poly}(W_1, W_2, 1/\epsilon')$ such that for all $n \geq M_1^{\epsilon'}$, both $|\mu_{n,i} - \mu_i| < \epsilon', \forall i \in A$ and
 725 $|T_{n,1}/n - \beta| \leq \epsilon'$ hold.

726 First we want to prove that for $n \geq M_1^{\epsilon'}$, if $O_n^{\epsilon/2}$ is nonempty, then P_n is nonempty. We prove
727 this by contradiction. Suppose P_n is empty. Then for all $i \neq 1$, $T_{n,i}/n \geq w_i^\beta$. Since $O_n^{\epsilon/2}$ is
728 nonempty, there exists some arm $\tilde{i} \neq 1$ such that $T_{n,\tilde{i}}/n > w_{\tilde{i}}^\beta + \epsilon/2$. In addition, for $n \geq M_1^{\epsilon'}$,
729 $T_{n,1}/n \geq \beta - \epsilon' > \beta - \epsilon/2$. Hence,

$$\begin{aligned} \sum_{i \in A} T_{n,i}/n &= T_{n,1}/n + T_{n,\tilde{i}}/n + \sum_{i \neq 1, \tilde{i}} T_{n,i}/n \\ &> \beta - \epsilon/2 + w_{\tilde{i}}^\beta + \epsilon/2 + \sum_{i \neq 1, \tilde{i}} w_i^\beta \\ &= \sum_{i \in A} w_i^\beta = 1, \end{aligned}$$

730 which leads to a contradiction since $\sum_{i \in A} T_{n,i}/n = (n-1)/n < 1$. Hence, for $n \geq M_1^{\epsilon'}$, if $O_n^{\epsilon/2}$ is
731 nonempty, then P_n is nonempty.

732 Next we will show that when n is sufficiently large, $I_n^{(2)} \notin O_n^{\epsilon/2}$. By Lemma 11, for $n \geq N_3$, we
733 have $I_n^{(1)} = I_n^* = 1$, and then for $i \neq 1$,

$$v_{n,i}^{(2)} = \sqrt{\sigma_{n,i}^2 + \sigma_{n,1}^2} f \left(\frac{\mu_{n,i} - \mu_{n,1}}{\sqrt{\sigma_{n,i}^2 + \sigma_{n,1}^2}} \right)$$

734 where $\sigma_{n,i}^2 = \sigma^2/T_{n,i}$ and $\sigma_{n,1}^2 = \sigma^2/T_{n,1}$. Note that for $n \geq M_1^{\epsilon'}$, $|\mu_{n,i} - \mu_i| < \epsilon'$, $\forall i \in A$ and
735 $|T_{n,1}/n - \beta| \leq \epsilon'$. Hence, for $n \geq \max\{N_3, M_1^{\epsilon'}\}$ and $i \in O_n^{\epsilon/2}$,

$$v_{n,i}^{(2)} < \sigma \left(\frac{1}{w_i^\beta + \epsilon/2} + \frac{1}{\beta - \epsilon'} \right)^{1/2} n^{-1/2} \phi \left(\frac{(\mu_i - \mu_1 + 2\epsilon')n^{1/2}}{\sigma \left[1/(w_i^\beta + \epsilon/2) + 1/(\beta - \epsilon') \right]^{1/2}} \right)$$

736 where the inequality uses Lemma 3. Note that $2\epsilon' < \epsilon < \Delta_{\min}/2$, so the value taken by $\phi(\cdot)$ is
737 negative. On the other hand, for $j \in P_n$,

$$\begin{aligned} v_{n,j}^{(2)} &> \sigma \left(\frac{1}{w_j^\beta} + \frac{1}{\beta + \epsilon'} \right)^{1/2} n^{-1/2} f \left(\frac{(\mu_j - \mu_1 - 2\epsilon')n^{1/2}}{\sigma \left[1/w_j^\beta + 1/(\beta + \epsilon') \right]^{1/2}} \right) \\ &> \sigma^4 \left(\frac{1}{w_j^\beta} + \frac{1}{\beta + \epsilon'} \right)^2 (-\mu_j + \mu_1 + 2\epsilon')^{-3} n^{-2} \phi \left(\frac{(\mu_j - \mu_1 - 2\epsilon')n^{1/2}}{\sigma \left[1/w_j^\beta + 1/(\beta + \epsilon') \right]^{1/2}} \right) \end{aligned}$$

738 where the last inequality is valid by Lemma 4 since there exists $M_2^{\epsilon'} = \text{poly}(1/\epsilon')$ such that for
739 $n \geq M_2^{\epsilon'}$, the value taken by both $f(\cdot)$ and $\phi(\cdot)$ is less than -2 . Let $M_3^{\epsilon'} \triangleq \max\{N_3, M_1^{\epsilon'}, M_2^{\epsilon'}\} =$
740 $\text{poly}(W_1, W_2, 1/\epsilon')$. For any $i, j \in A$ such that $i \neq j$ and $i, j \neq 1$, we define the following constant
741 in terms of ϵ

$$C_{i,j}^\epsilon \triangleq \frac{(\mu_i - \mu_1)^2}{1/(w_i^\beta + \epsilon/2) + 1/\beta} - \frac{(\mu_j - \mu_1)^2}{1/w_j^\beta + 1/\beta},$$

742 and we let

$$C_{\min}^\epsilon \triangleq \min_{\substack{i \neq j \\ i, j \neq 1}} C_{i,j}^\epsilon,$$

743 and for $\epsilon' \in (0, \epsilon/2)$, we define the following function of ϵ'

$$g_{i,j}^\epsilon(\epsilon') \triangleq \frac{(\mu_i - \mu_1 + 2\epsilon')^2}{1/(w_i^\beta + \epsilon/2) + 1/(\beta - \epsilon')} - \frac{(\mu_j - \mu_1 - 2\epsilon')^2}{1/w_j^\beta + 1/(\beta + \epsilon')}.$$

744 We know that

$$\frac{(\mu_2 - \mu_1)^2}{1/w_2^\beta + 1/\beta} = \dots = \frac{(\mu_k - \mu_1)^2}{1/w_k^\beta + 1/\beta},$$

745 so each $C_{i,j}^\epsilon > 0$, and thus $C_{\min}^\epsilon > 0$. Since each $g_{i,j}^\epsilon(\epsilon')$ is increasing as ϵ' is decreasing to 0,
 746 and $\lim_{\epsilon' \rightarrow 0} g_{i,j}^\epsilon(\epsilon') = C_{i,j}^\epsilon \geq C_{\min}^\epsilon$, there exists a threshold $\epsilon_{i,j} = \text{poly}(\epsilon) \in (0, \epsilon/2)$ such that
 747 $g_{i,j}^\epsilon(\epsilon_{i,j}) \geq C_{\min}^\epsilon/2$ (note that $\epsilon < 1$). We let

$$\epsilon_{\min} \triangleq \min_{\substack{i \neq j \\ i, j \neq 1}} \epsilon_{i,j}.$$

748 Then for $n \geq M_3^{\epsilon_{\min}}$, for all $i \in O_n^{\epsilon/2}$ and $j \in P_n$,

$$\frac{v_{n,j}^{(2)}}{v_{n,i}^{(2)}} > D_{i,j}^\epsilon n^{-3/2} \exp\left(\frac{C_{\min}^\epsilon n}{4\sigma^2}\right) \geq D_{\min}^\epsilon n^{-3/2} \exp\left(\frac{C_{\min}^\epsilon n}{4\sigma^2}\right), \quad (12)$$

749 where

$$D_{i,j}^\epsilon \triangleq \frac{\sigma^4 \left(\frac{1}{w_j^\beta} + \frac{1}{\beta + \epsilon_{\min}}\right)^2 (-\mu_j + \mu_1 + 2\epsilon_{\min})^{-3}}{\sigma \left(\frac{1}{w_i^\beta + \epsilon/2} + \frac{1}{\beta - \epsilon_{\min}}\right)^{1/2}}$$

750 and

$$D_{\min}^\epsilon \triangleq \min_{\substack{i \neq j \\ i, j \neq 1}} D_{i,j}^\epsilon.$$

751 Since $\epsilon_{\min} = \text{poly}(\epsilon)$, there exists $M_4^\epsilon = \text{poly}(1/\epsilon, \epsilon)$ such that for $n \geq M_4^\epsilon$, the right hand side
 752 of (12) is greater than 1. Hence, for $n \geq M_5^\epsilon \triangleq \max\{M_3^{\epsilon_{\min}}, M_4^\epsilon\}$ where $\epsilon_{\min} = \text{poly}(\epsilon)$, we
 753 have $v_{n,j}^{(2)} > v_{n,i}^{(2)}$ for all $i \in O_n^{\epsilon/2}$ and $j \in P_n$, which implies $I_n^{(2)} \notin O_n^{\epsilon/2}$. Note that $M_5^\epsilon =$
 754 $\text{poly}(W_1, W_2, 1/\epsilon, \epsilon)$.

755 Finally we will prove when n is sufficiently large, O_n^ϵ is empty. Let $M^\epsilon \triangleq \max\{M_5^\epsilon, 2/\epsilon\}$. There
 756 are two following cases on the set $O_{M^\epsilon}^{\epsilon/2}$.

757 1. $|O_{M^\epsilon}^{\epsilon/2}| = 0$

758 We will prove by induction that for all $n \geq M^\epsilon$, O_n^ϵ is empty. For $n = M^\epsilon$, O_n^ϵ is empty
 759 since $O_n^\epsilon \subseteq O_n^{\epsilon/2}$ and $O_n^{\epsilon/2}$ is empty. Now we suppose that O_n^ϵ is empty for some $n \geq M^\epsilon$,
 760 and we want to show that O_{n+1}^ϵ is empty.

761 Note that O_n^ϵ is empty, and then only $I_n^{(1)}$ and $I_n^{(2)}$ may enter O_{n+1}^ϵ . We known that for
 762 $n \geq M^\epsilon$, $I_n^{(1)} = 1$, which implies that $I_n^{(2)} \neq 1$ and only $I_n^{(2)}$ may enter O_{n+1}^ϵ . In addition,
 763 for $n \geq M^\epsilon$, we have proved that $I_n^{(2)} \notin O_n^{\epsilon/2}$, which implies $T_{n, I_n^{(2)}}/n - w_{I_n^{(2)}}^\beta \leq \epsilon/2$.

764 Since $n \geq M^\epsilon \geq 2/\epsilon$, $T_{n+1, I_n^{(2)}}/(n+1) - w_{I_n^{(2)}}^\beta \leq (T_{n, I_n^{(2)}} + 1)/n - w_{I_n^{(2)}}^\beta \leq 1/n + \epsilon/2 \leq \epsilon$,
 765 which implies $I_n^{(2)} \notin O_{n+1}^\epsilon$, i.e., $I_n^{(2)}$ will not enter O_{n+1}^ϵ . Hence, if O_n^ϵ is empty, then
 766 O_{n+1}^ϵ is empty.

767 Therefore, by induction, for all $n \geq M^\epsilon$, O_n^ϵ is empty.

768 2. $|O_{M^\epsilon}^{\epsilon/2}| \geq 1$

769 Similarly to the proof for case 1, we can show that for any arm $i \notin O_{M^\epsilon}^{\epsilon/2}$, it will not enter
 770 any O_n^ϵ for $n \geq M^\epsilon$.

771 Now let us consider arm $i \in O_{M^\epsilon}^{\epsilon/2}$. Let L_i^ϵ be the time such that $i \in O_n^{\epsilon/2}$ for $n \in$
 772 $[M^\epsilon, L_i^\epsilon - 1]$ and $i \notin O_{L_i^\epsilon}^{\epsilon/2}$. Similar to the proof for case 1, we can prove that for i will not
 773 enter any O_n^ϵ for $n \geq L_i^\epsilon$.

774 Let $M_6^\epsilon \triangleq \max_{i \in O_{M^\epsilon}^{\epsilon/2}} L_i^\epsilon$. For $n \geq M_6^\epsilon$, O_n^ϵ is empty. Note that $M_6^\epsilon =$
 775 $\text{poly}(W_1, W_2, 1/\epsilon, \epsilon)$.

776 Combining the above two cases, we conclude that there exists $N_6^\epsilon = \text{poly}(W_1, W_2, 1/\epsilon, \epsilon)$ such that
 777 for all $n \geq N_6^\epsilon$, O_n^ϵ is empty.

778 □

779 Based on Lemma 14, we can easily prove that when n is large, the under-sampled set is also empty,
 780 which immediately establishes Proposition 4.

781 **Proof of Proposition 4.** Given $\epsilon > 0$, by Lemmas 13 and 14, there exists $M_1^{\epsilon/k} =$
 782 $\text{poly}(W_1, W_2, 1/\epsilon, \epsilon)$ such that for $n \geq M_1^{\epsilon/k}$, $|T_{n,1}/n - w_1^\beta| \leq \epsilon/k$ where $w_1^\beta = \beta$ and
 783 $T_{n,i}/n - w_i^\beta \leq \epsilon/k$ for all $i \in A \setminus \{1\}$. Suppose there exists $i' \in A$ such that $T_{n,i'}/n - w_{i'}^\beta < -\epsilon$.
 784 Then

$$\begin{aligned} \sum_{i \in A} T_{n,i}/n &= T_{n,i'}/n + \sum_{i \neq i'} T_{n,i}/n \\ &< w_{i'}^\beta - \epsilon + \sum_{i \neq i'} (w_i^\beta + \epsilon/k) \\ &= \sum_{i \in A} w_i^\beta + [-\epsilon + (k-1)\epsilon/k] \\ &= 1 - \epsilon/k. \end{aligned}$$

785 On the other hand, for $n \geq k/\epsilon$, $\sum_{i \in A} T_{n,i}/n = (n-1)/n \geq 1 - \epsilon/k$, which leads to a contradiction.
 786 Hence, for $n \geq N_7^\epsilon = \max\{M_1^{\epsilon/k}, k/\epsilon\}$, for all $i \in A$, we have $-\epsilon \leq T_{n,i}/n - w_i^\beta \leq \epsilon/k$, which
 787 leads to $|T_{n,i}/n - w_i^\beta| < \epsilon$. Note that $N_7^\epsilon = \text{poly}(W_1, W_2, 1/\epsilon, \epsilon)$.

788 F.5 Proof of Theorem 1

789 For any $\epsilon > 0$, by Propositions 3 and 4, for $n \geq N_\beta^\epsilon \triangleq \{N_2^\epsilon, N_7^\epsilon\}$, we have

$$|\mu_{n,i} - \mu_i| \leq \epsilon \quad \text{and} \quad |T_{n,i}/n - w_i^\beta| \leq \epsilon \quad \forall i \in A.$$

790 Note that $N_\beta^\epsilon = \text{poly}(W_1, W_2, 1/\epsilon, \epsilon)$. By Lemmas 5 and 6, we have $\mathbb{E}[e^{\lambda W_1}] < \infty$ and $\mathbb{E}[e^{\lambda W_2}] <$
 791 ∞ for all $\lambda > 0$, which implies that the expected value of any polynomial of W_1 and W_2 is finite,
 792 and thus $\mathbb{E}[N_\beta^\epsilon] < \infty$. By definition, $T_\beta^\epsilon \leq N_\beta^\epsilon$, so $\mathbb{E}[T_\beta^\epsilon] \leq \mathbb{E}[N_\beta^\epsilon] < \infty$.

793 Since ϵ can be arbitrary small, for any sample path (up to a set of measure zero), we have

$$\lim_{n \rightarrow \infty} \frac{T_{n,i}}{n} = w_i^\beta \quad \forall i \in A.$$